

CS-570

Statistical Signal Processing

Lecture 10: Quantization & sampling

Spring Semester 2019

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Matrix Completion (MC)

Relaxation

$$\min\{ \|\mathbf{M}\|_* : \mathcal{A}(\mathbf{X}) = \mathcal{A}(\mathbf{M}) \}$$

Performance

$$\left\| M - M^* \right\|_F^2 \leq 4 \sqrt{\frac{(2+p)\min(n_1, n_2)}{p}} \delta + 2\delta,$$

where p = fraction of known entries = $\frac{m}{n_1 n_2} = \frac{|\Omega|}{n_1 n_2}$

Noisy case

$$\min\{ \|\mathbf{M}\|_* : \|\mathcal{A}(\mathbf{X}) - \mathcal{A}(\mathbf{M})\|_F^2 \leq \epsilon \}$$



CS and MC

	<i>Sparse recovery</i>	<i>Rank minimization</i>
Unknown	Vector x	Matrix A
Observations	$y = Ax$	$y = L[A]$ (linear map)
Combinatorial objective	$\#\{\mathbf{x}_i \neq 0\} = \ \mathbf{x}\ _0$	$\text{rank}(A) = \#\{\sigma_i(A) \neq 0\} = \ \sigma(A)\ _0$
Convex relaxation	$\ \mathbf{x}\ _1 = \sum_i \mathbf{x}_i $	$\ A\ _* = \sum_i \sigma_i(A)$
Algorithmic tools	Linear programming	Semidefinite programming

Yi Ma et al, "Matrix Extensions to Sparse Recovery", CVPR2009



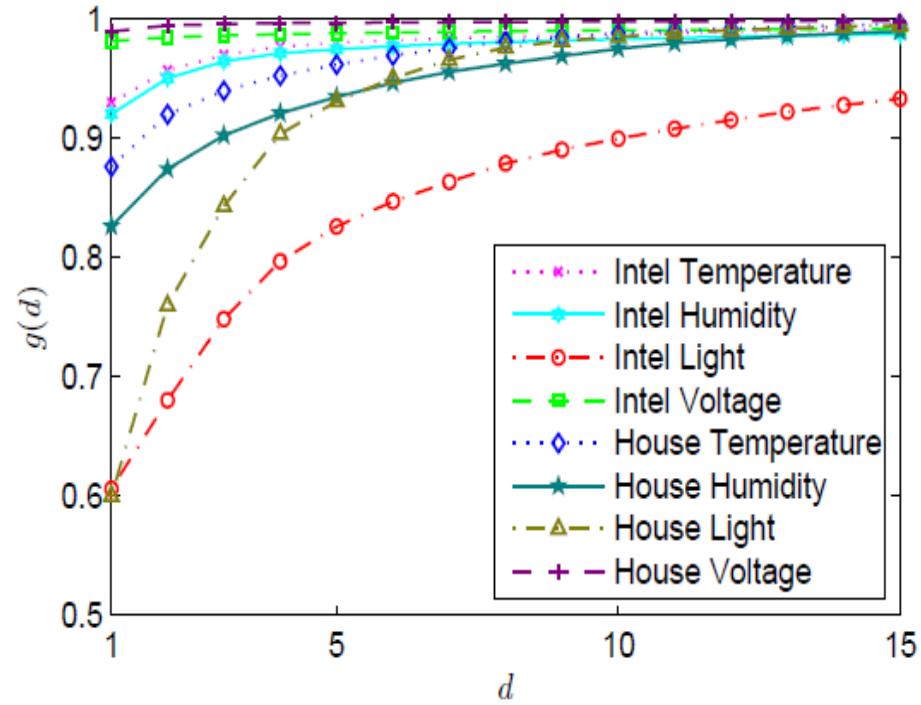
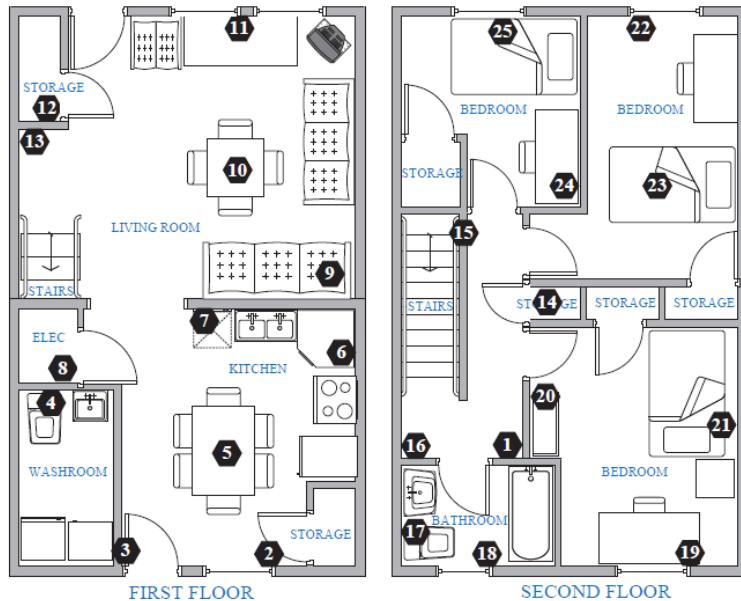
Applications of MC

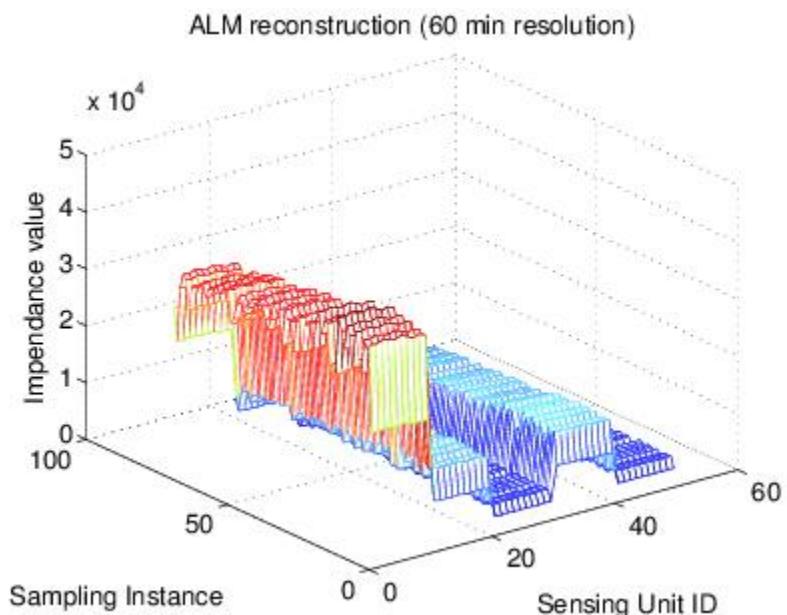
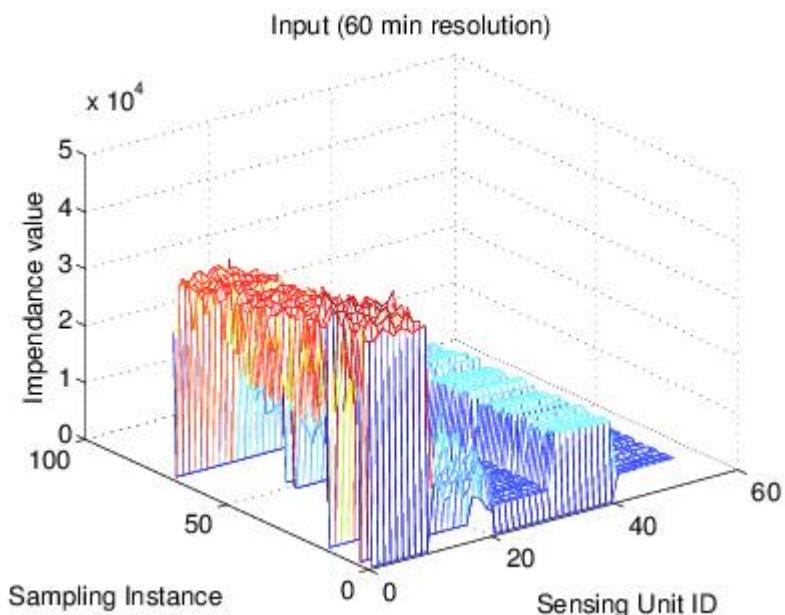
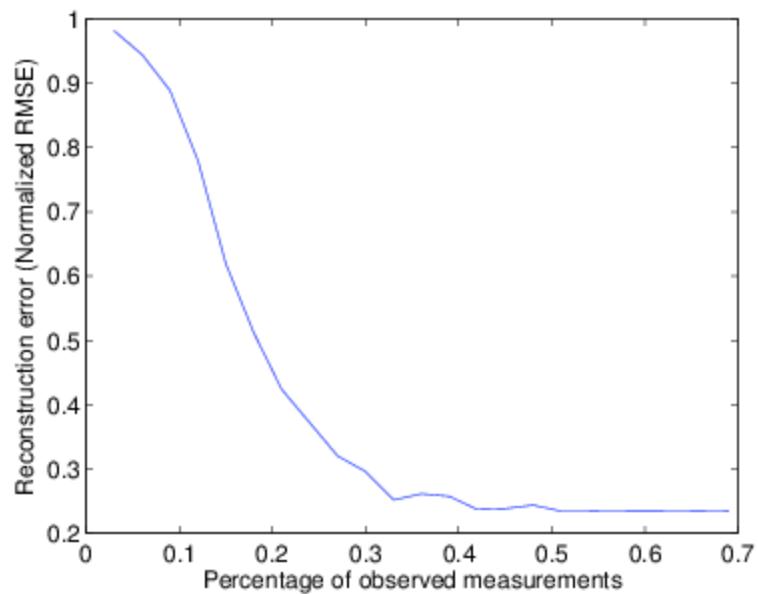
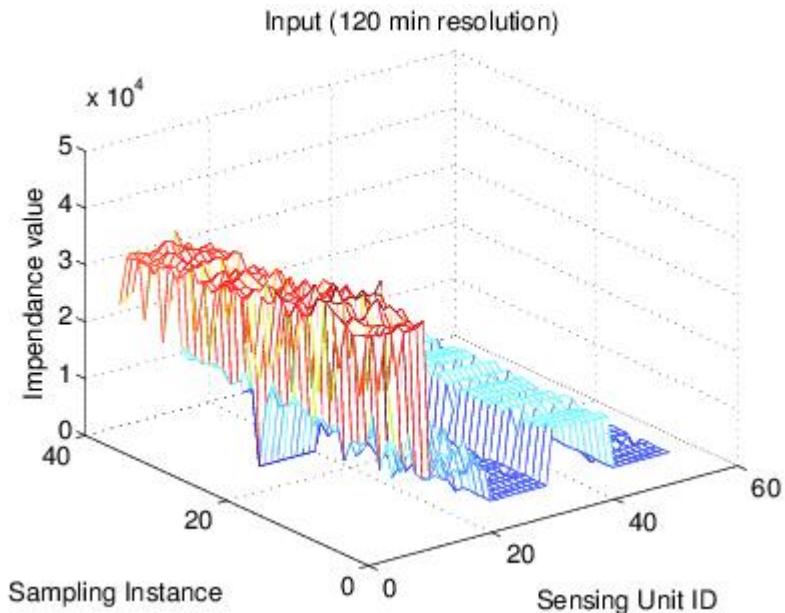
- Recommendation systems
 - Matrix (user, preference/quality/intention)
- Sensor localization
 - Matrix (location, physical quantity)
- Data recovery in Wireless Sensor Networks
 - Matrix (sensor, time)



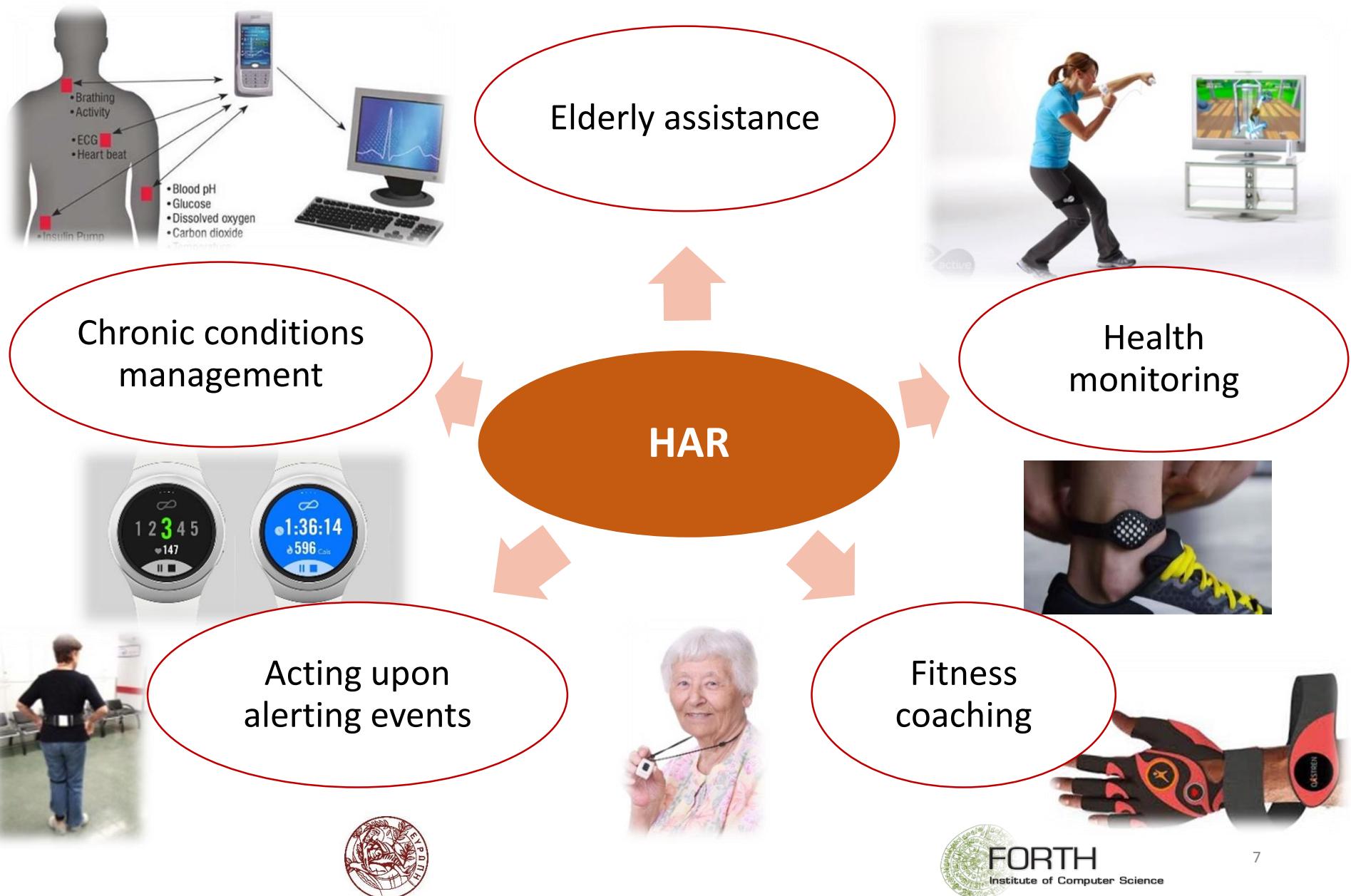
Data Gathering

- STCDG: An Efficient Data Gathering Algorithm Based on Matrix Completion for Wireless Sensor Networks





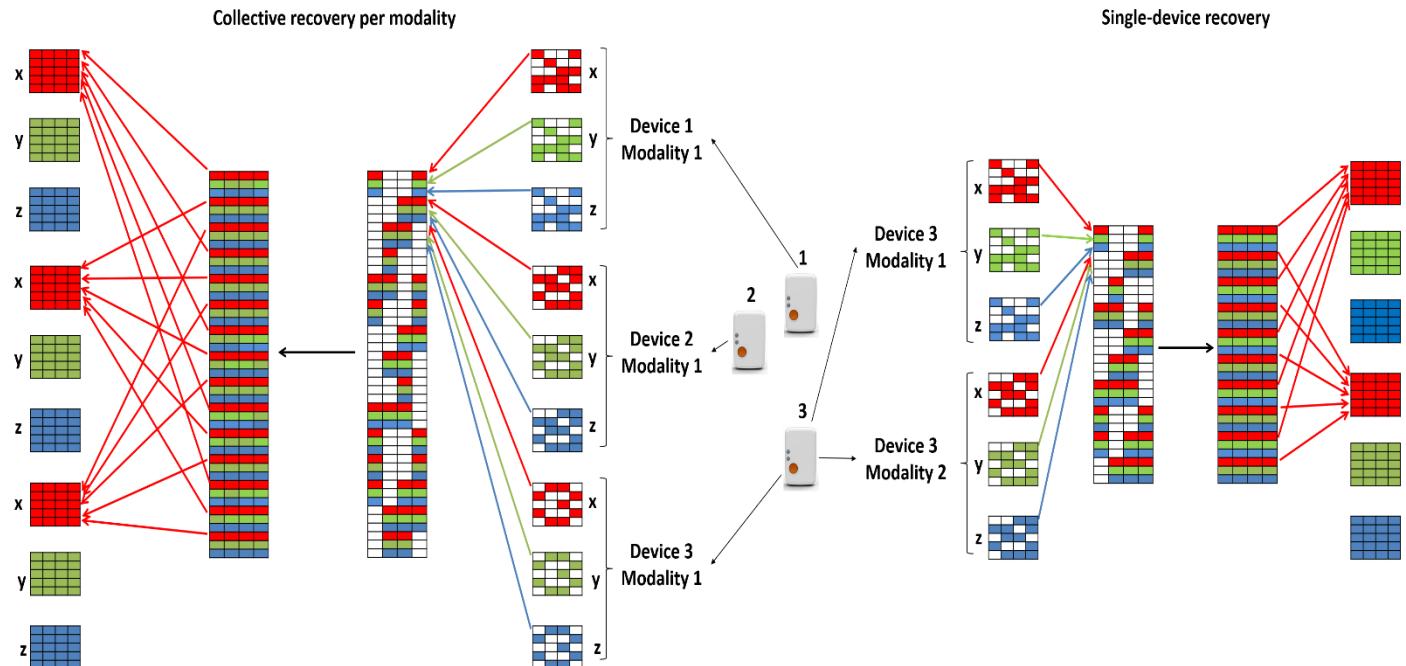
WSNs for Human Activity Recognition



Single-device vs collective recovery: matrices

Scenario 2 Collective per modality

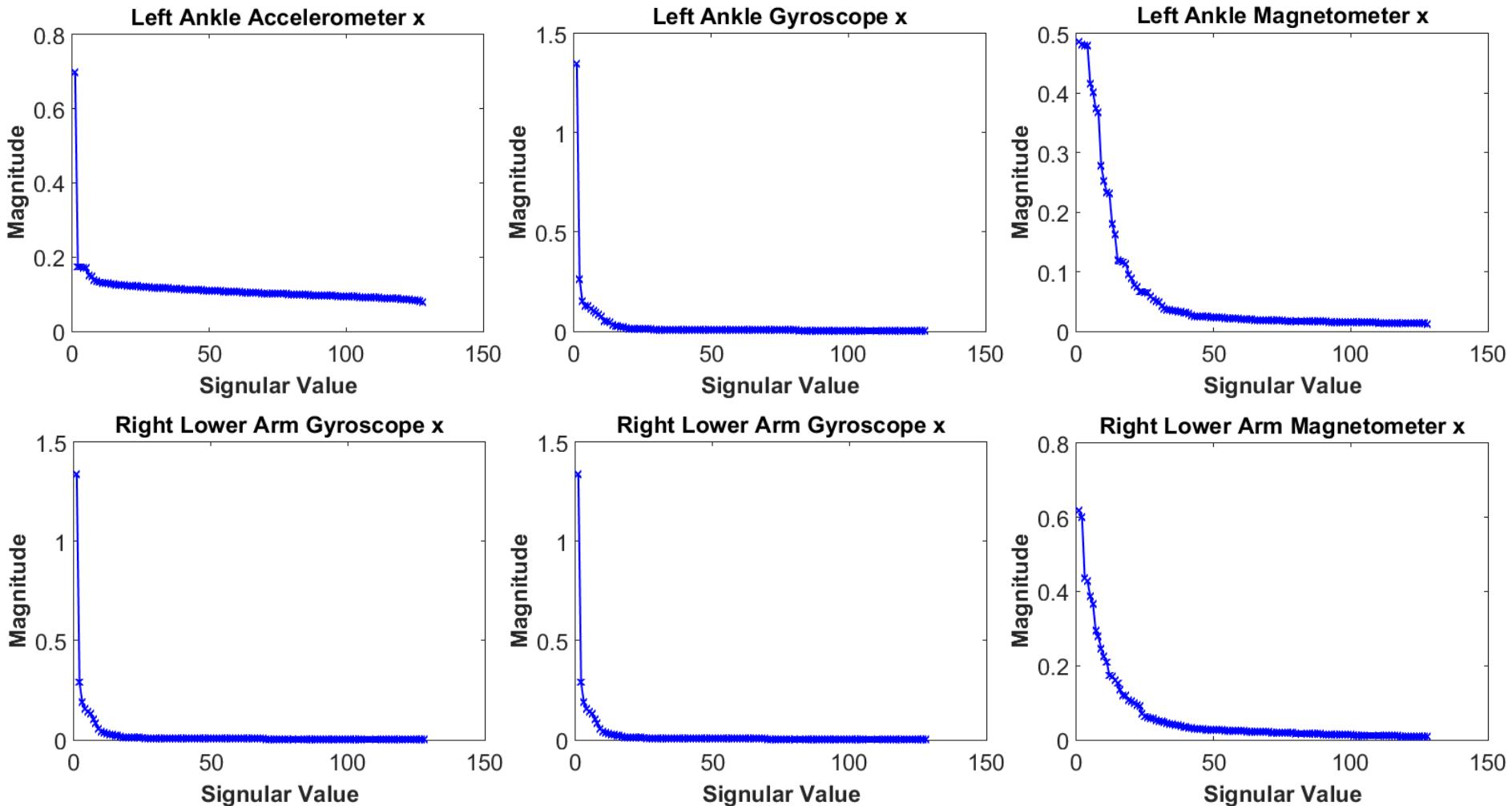
Scenario 1 Single-device



Scenario 3: Overall collective recovery structured similarly



Body sensor network



RTT estimation

Decentralized Matrix Factorization by Stochastic Gradient Descent (DMFSGD),

Estimation of end-to-end network distances

- Network nodes exchange messages with each other
- Each node collects and processes local measurements

$$D \approx X \begin{matrix} \text{r columns} \\ \times \\ Y^T \end{matrix} = \hat{D}$$

The diagram illustrates the matrix factorization process. On the left, a square matrix D is shown with various numerical values in its entries. Above it, the symbol \approx indicates that D is approximately equal to the product of two matrices X and Y^T . The matrix X has r columns, as indicated by a bracket above it. To the right of the \times operator, the transpose of Y , denoted as Y^T , is shown. The result of the multiplication is a matrix \hat{D} , which also contains numerical values. The diagonal entries of both D and \hat{D} are empty.

Fig. 2. Network distance prediction by matrix factorization. Note that the diagonal entries of D and \hat{D} are empty.

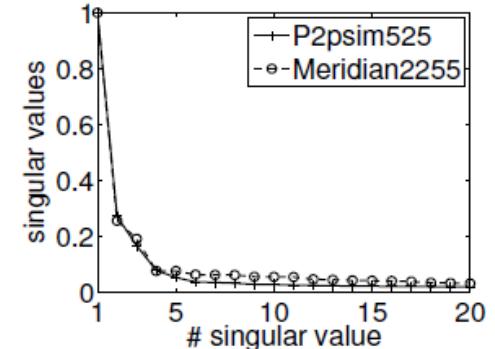
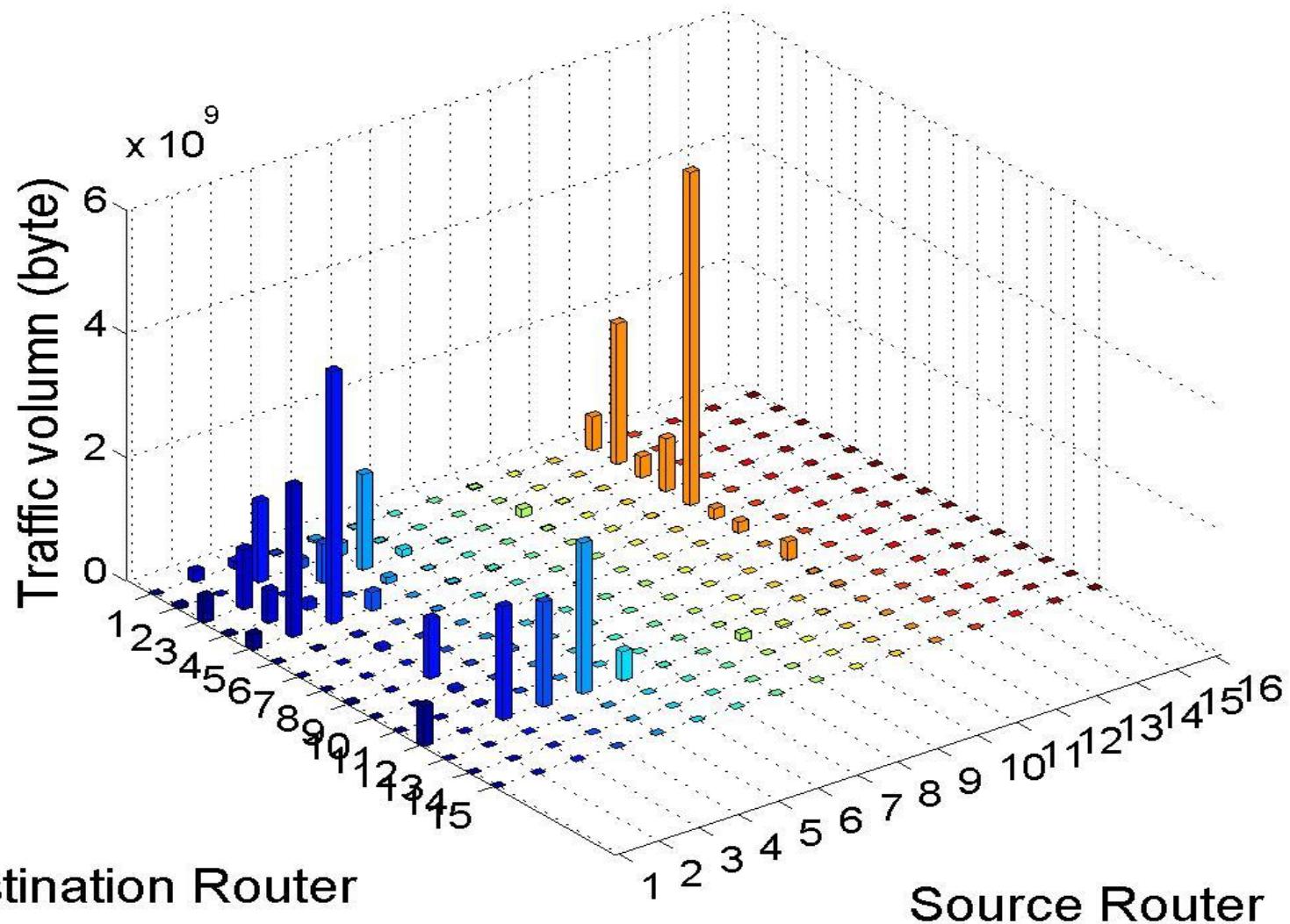
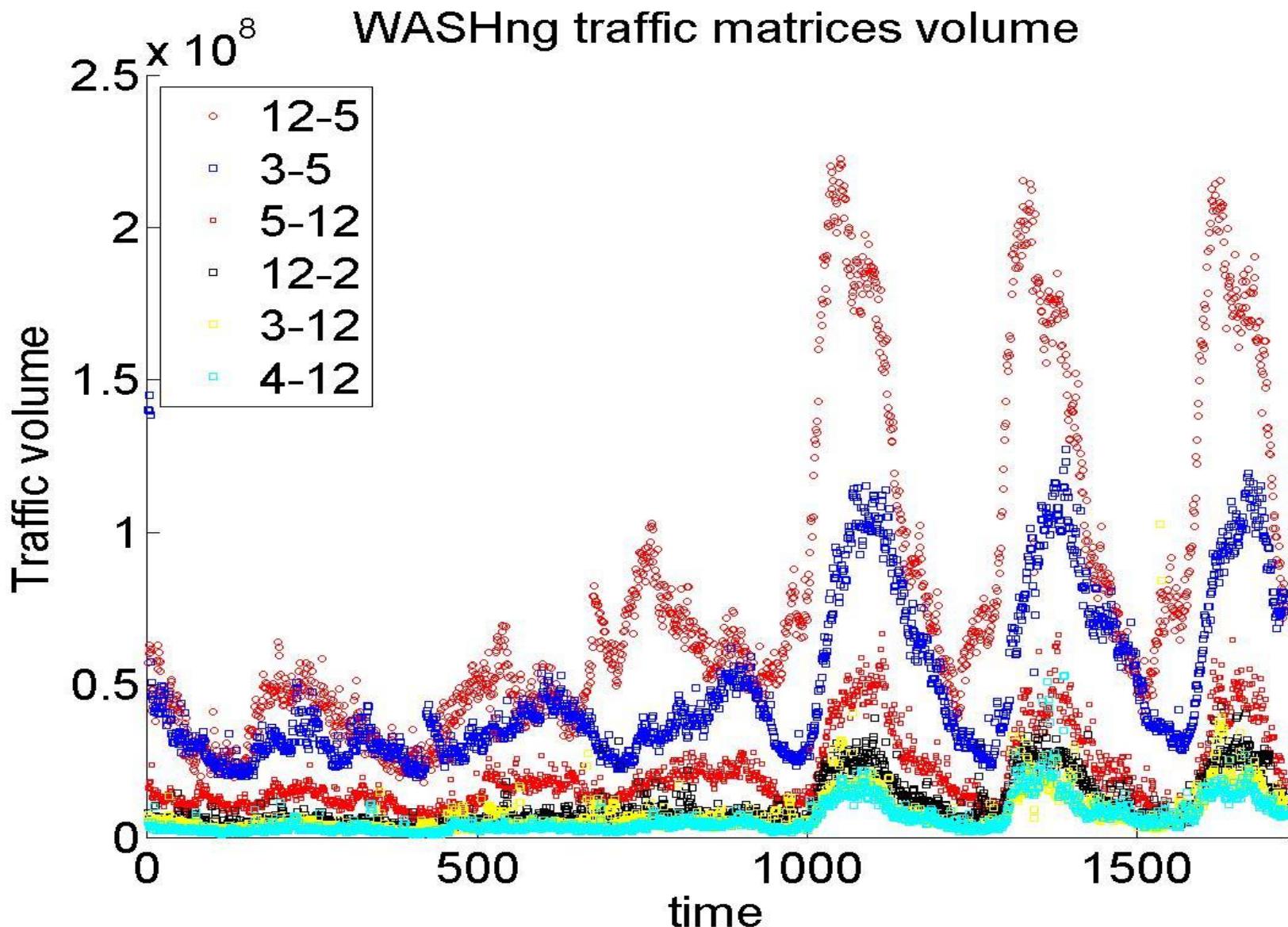


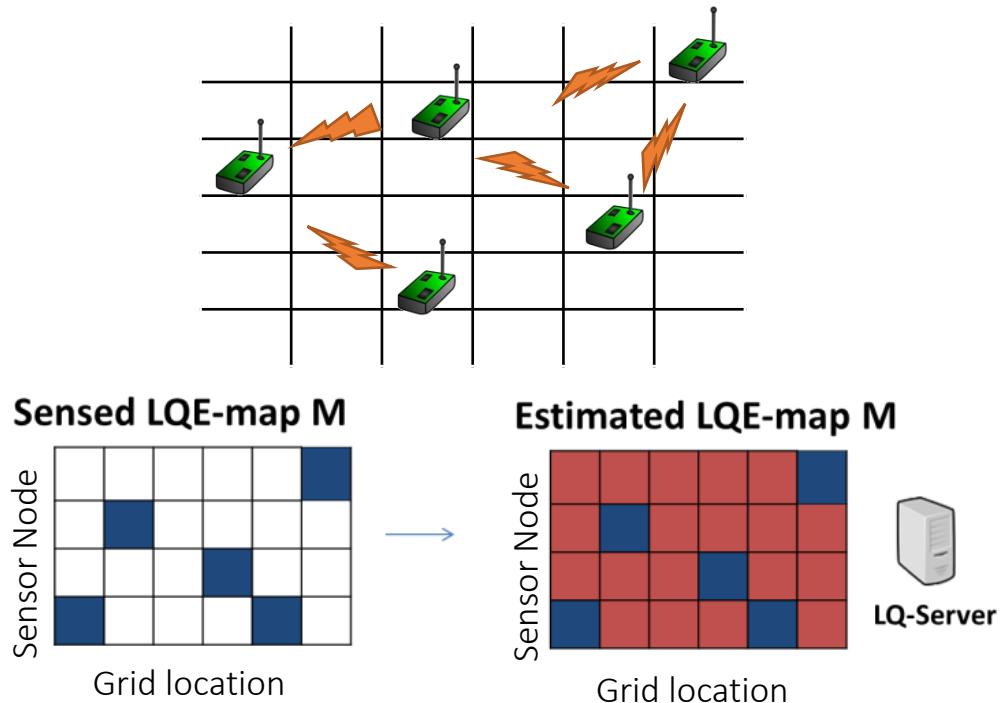
Fig. 3. The singular values of a RTT matrix of 2255×2255 , extracted from the Meridian dataset [30] and called “Meridian2255”, and of a RTT matrix of 525×525 , extracted from the P2psim dataset [30] and called “P2psim525”. The singular values are normalized so that the largest singular values of both matrices are equal to 1.

Traffic Matrix of router WASHng

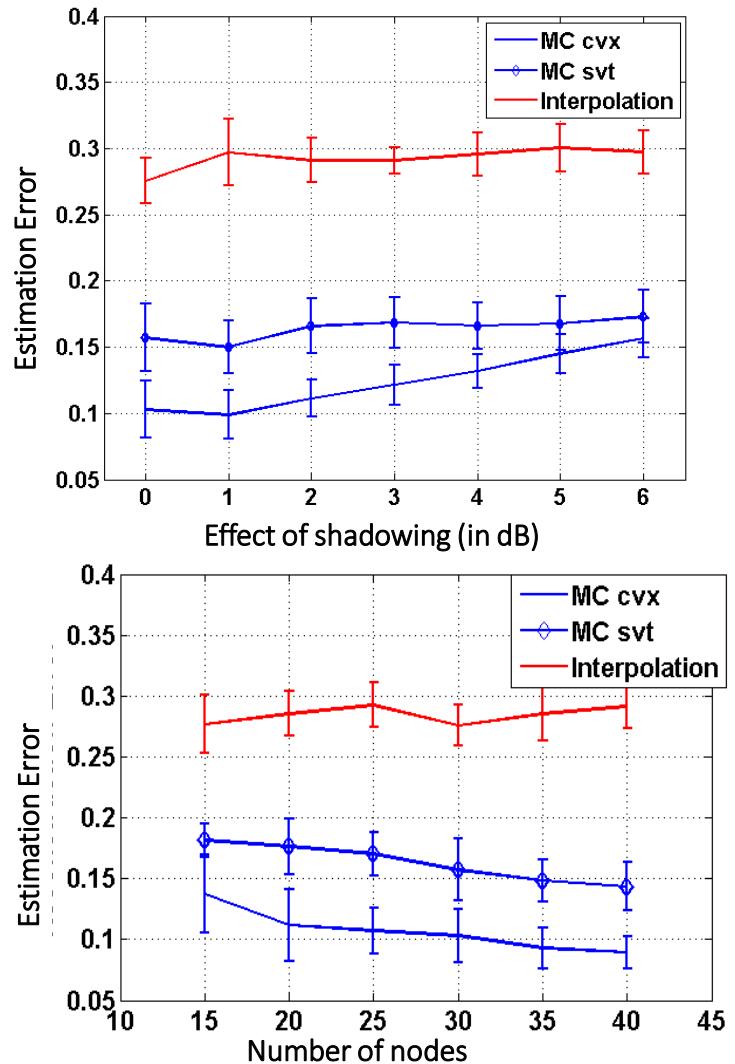




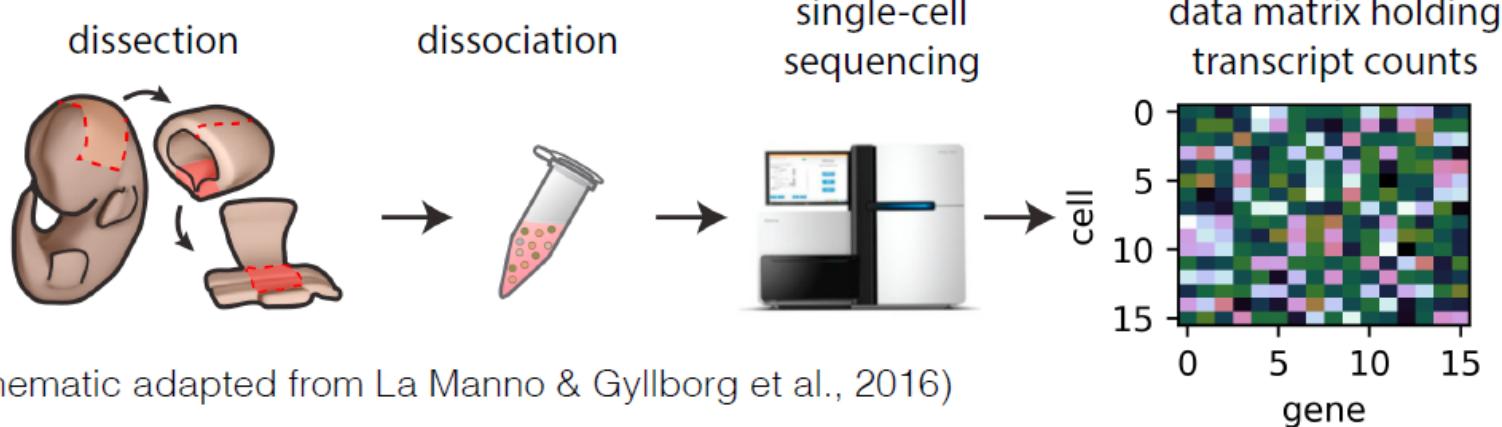
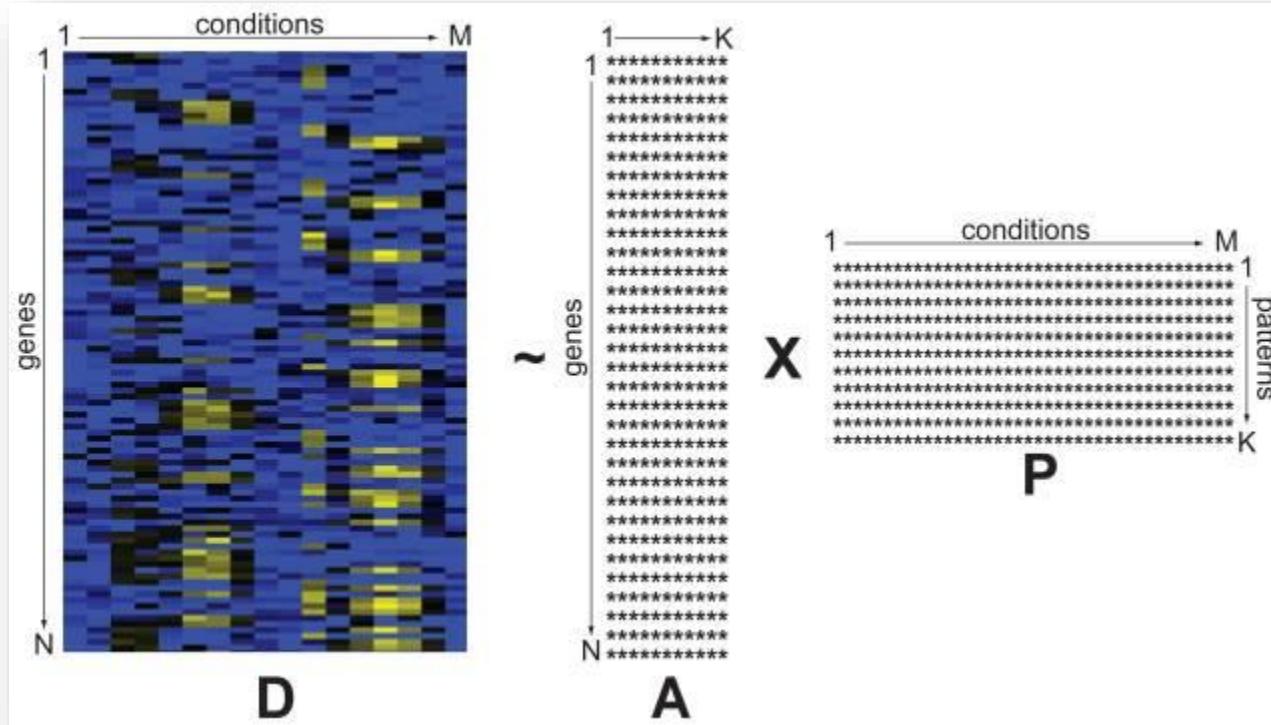
LQM Estimation



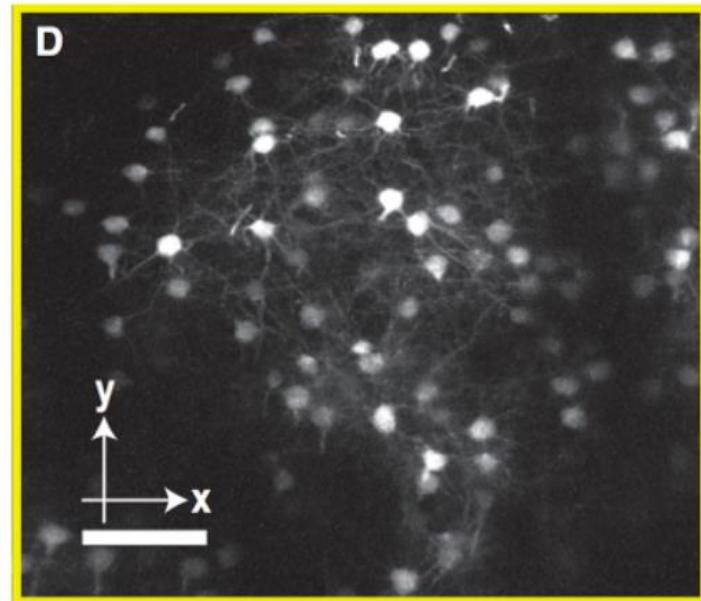
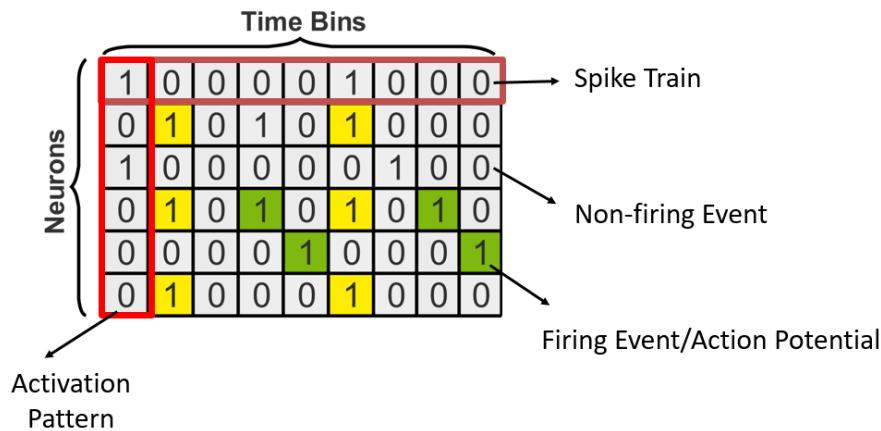
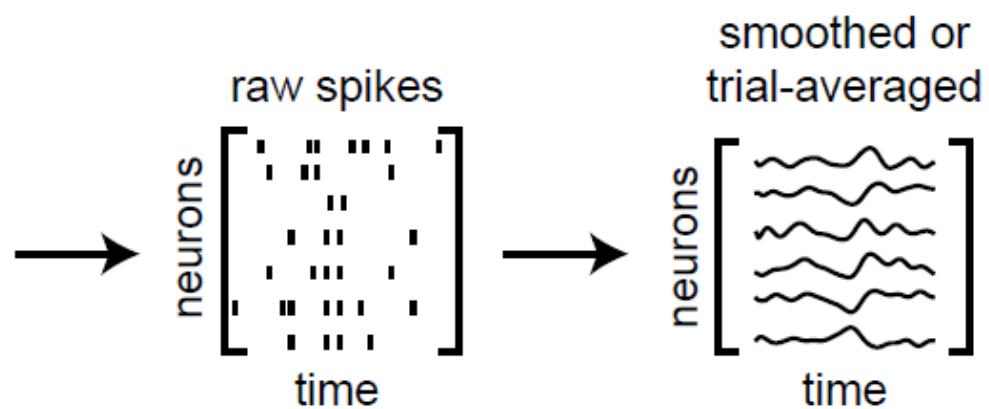
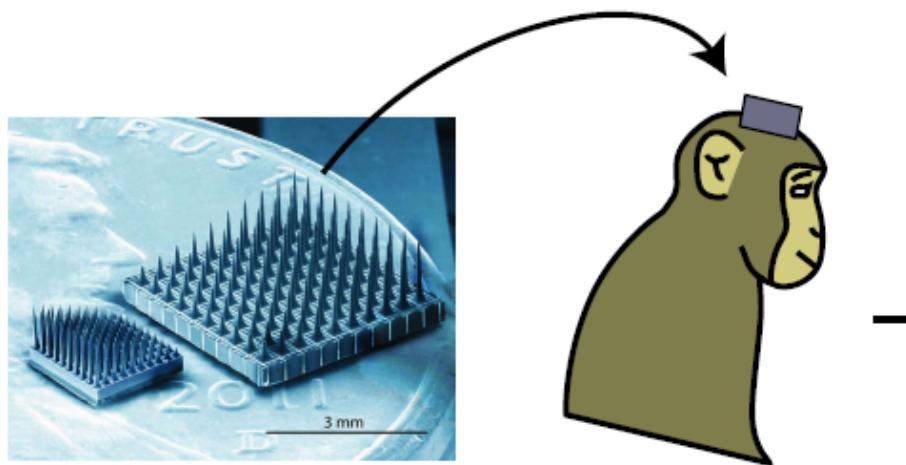
- Dataset
 - Testbed @ FORTH (144m², 1x1m grid)
 - RSSI values (channel quality)
 - 13 IEEE802.11b/g channels



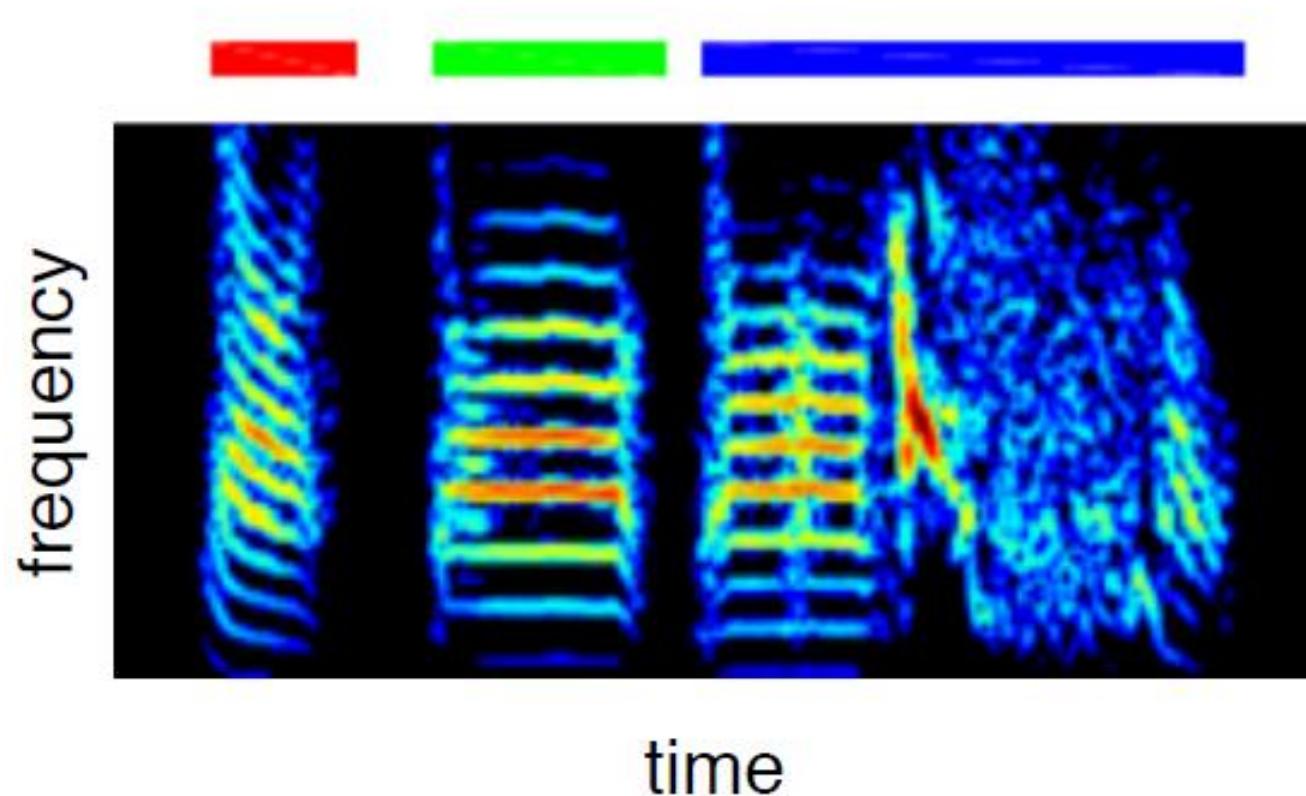
Microarray data



Neural Activity



Audio



Transductive Classification

- Unknown labels and corrupted or missing training features and corresponding labels leading to incomplete matrix

$$\mathbf{Z} = \begin{bmatrix} \mathbf{Y}_{tr} & \mathbf{Y}_{tst} \\ \mathbf{X}_{tr} & \mathbf{X}_{tst} \end{bmatrix} \Rightarrow \begin{bmatrix} \mathbf{Y}_{tr} & ? \\ \mathbf{X}_{tr} & \mathbf{X}_{tst} \end{bmatrix}$$

Labels
Features

Train | Test

- Formulate classification as rank minimization

$$\begin{array}{ll} \min & \mu \|\mathbf{Z}\|_* + l_X(\mathbf{Z}_X) + \lambda l_Y(\mathbf{Z}_Y) \\ \text{s.t.} & \mathbf{Z}_X \geq \mathbf{0} \text{ and } \mathbf{Z}_1 = \mathbf{1}^T \end{array}$$

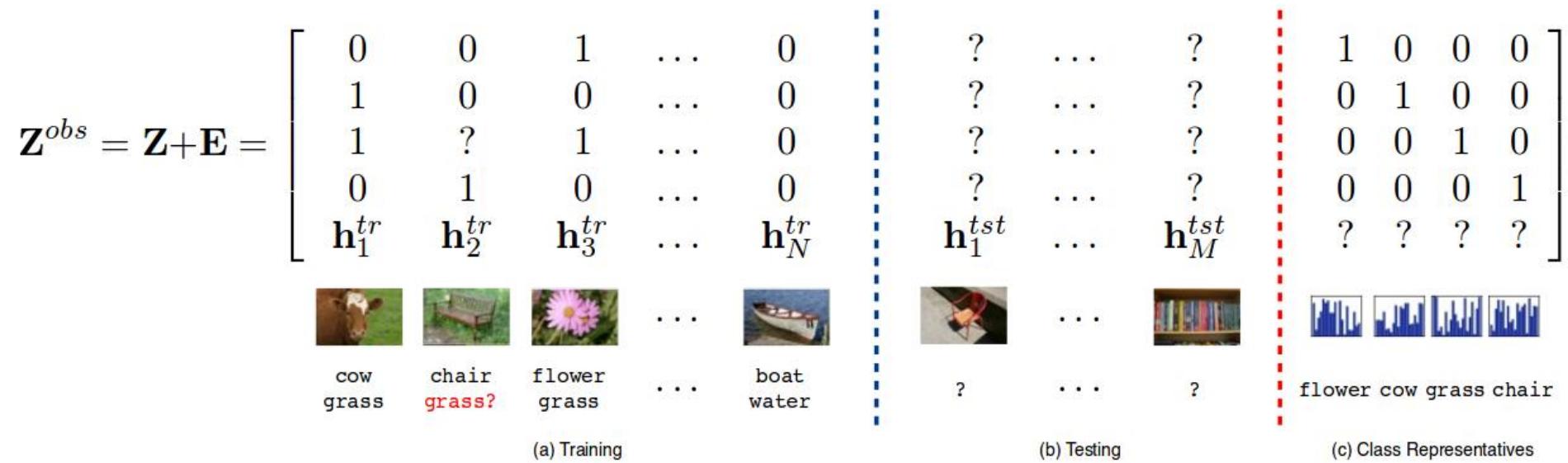
where:

$$l_X(\mathbf{Z}_X) = \sum_{ij \in \mathbf{Z}_X} (z_{ij} - z_{0ij})^2$$

$$l_Y(\mathbf{Z}_Y) = \sum_{ij \in \mathbf{Y}_{train}} \frac{1}{\gamma} \log(1 + \exp(-\gamma z_{ij} z_{0ij}))$$



Multi-label image classification

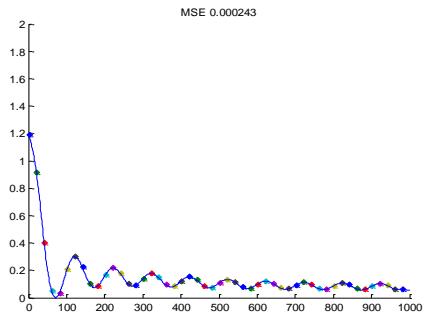


Analog-to-digital transformation

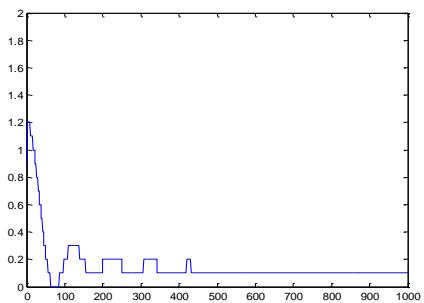


Signal Sensing

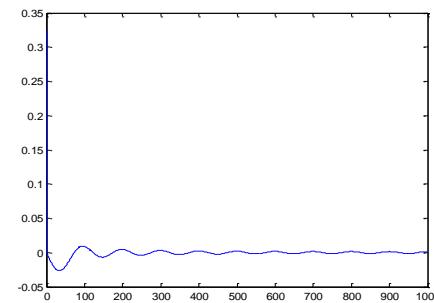
Sample



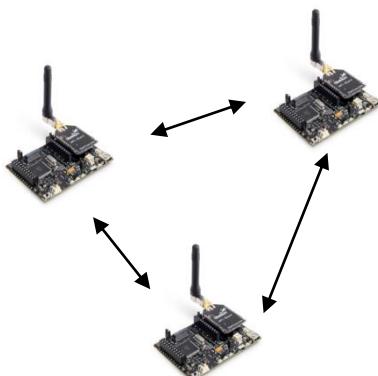
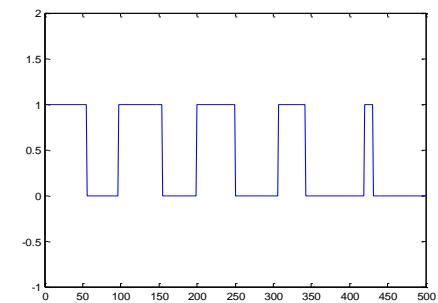
Quantize



Process



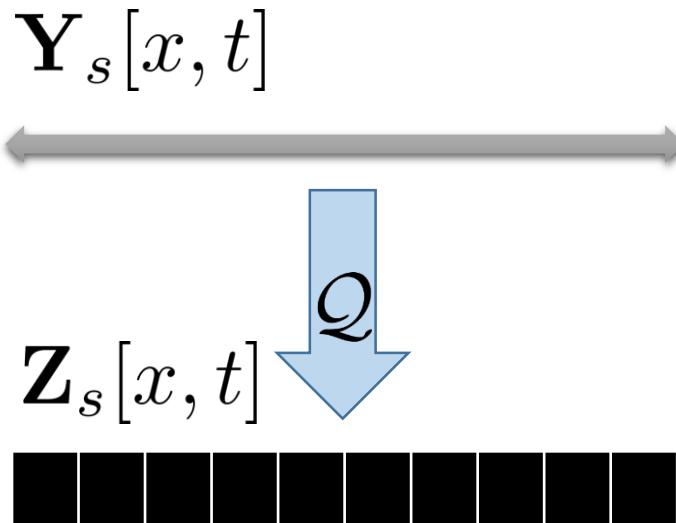
Transmit



Signal Quantization

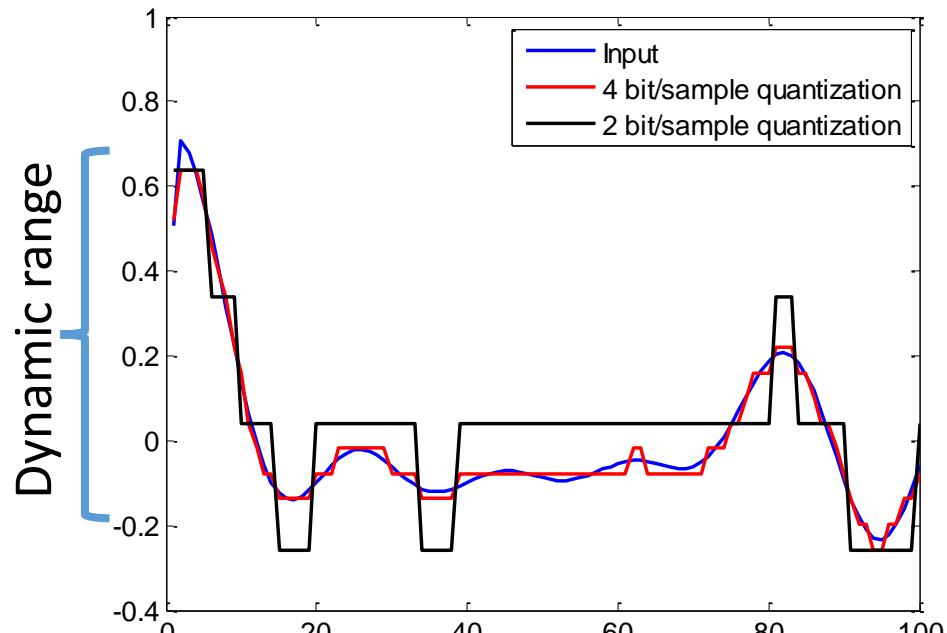
Map values to symbols

- Bits per measurements



$$|\mathcal{C}| = \mathcal{B} = 2^R$$

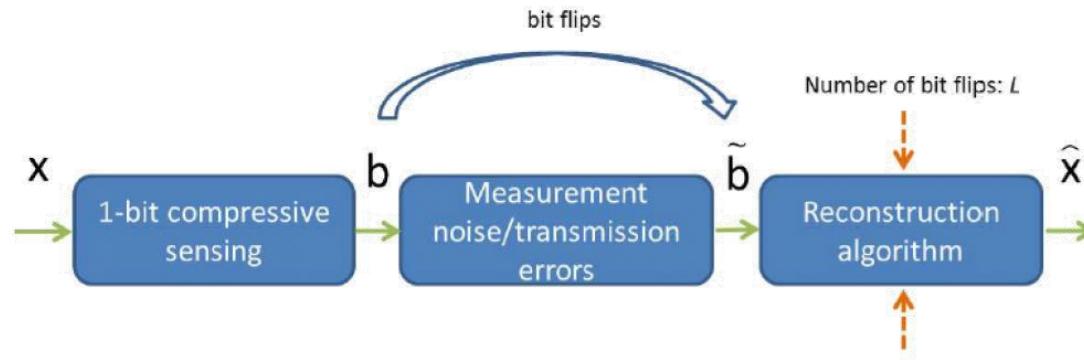
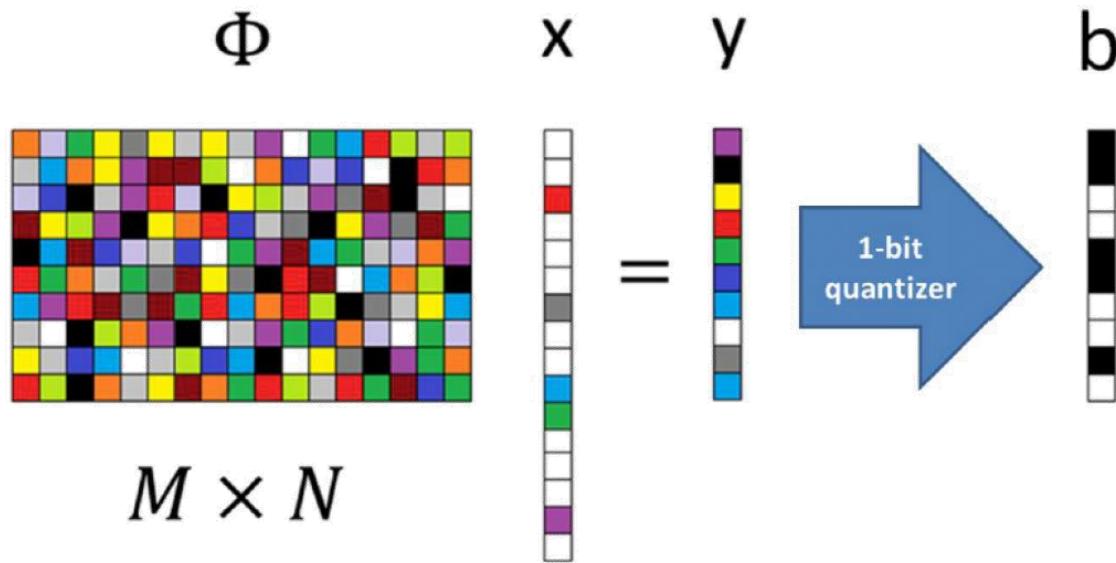
$$Q(x) = \Delta \cdot \left\lfloor \frac{x}{\Delta} + \frac{1}{2} \right\rfloor = \Delta \cdot \text{floor} \left(\frac{x}{\Delta} + \frac{1}{2} \right)$$



$$\mathcal{T} = \{t_1, \dots, t_{\mathcal{B}} | t_i < t_j, \forall i, j\}$$



1-Bit Compressive Sensing



End-to-End CS imaging

- Difference aspects of sensing process

- Dictionary

- Noise sources

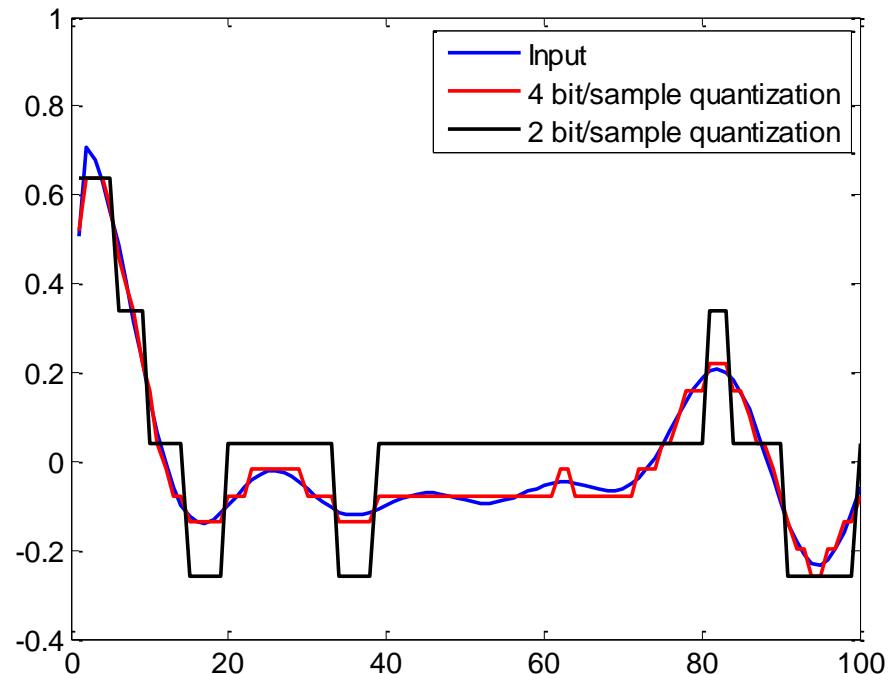
• Quantization

- ADC

- Storage

- Transmission

- Entropy Coding



(Scalar) Signal Quantization

Map from real to indexed $\mathcal{Q} : \mathbb{R} \rightarrow 2^R$

➤ Codebook

$$|\mathcal{C}| = \mathcal{B} = 2^R$$

➤ Thresholds

$$\mathcal{T} = \{t_1, \dots t_{\mathcal{B}} | t_i < t_j, \forall i, j\}$$

➤ Dynamic range

$$M = |\max(x) - \min(x)|$$

➤ Saturation

$$\hat{f} = \min(f, M)$$

Uniform scalar

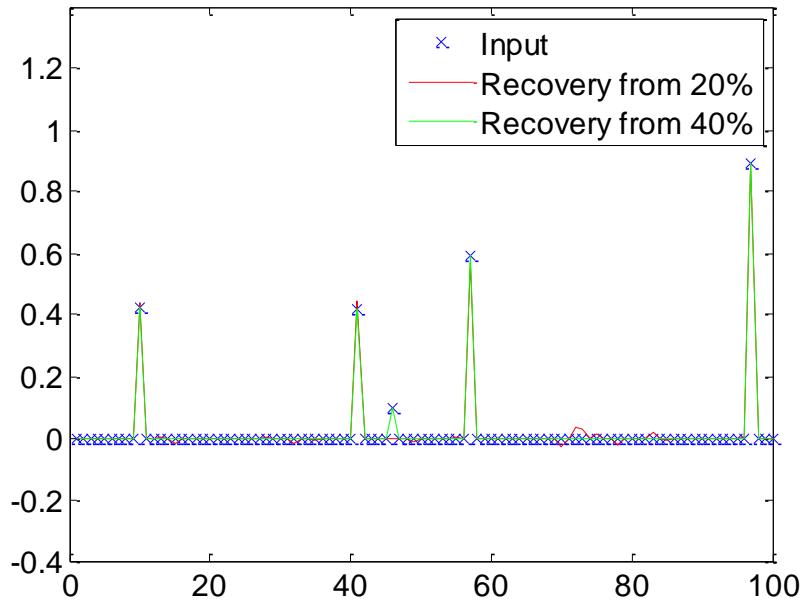
Non-uniform scalar: Lloyd-Max conditions

$$|t_j - t_i| = \Delta \quad n = \log_2 M$$

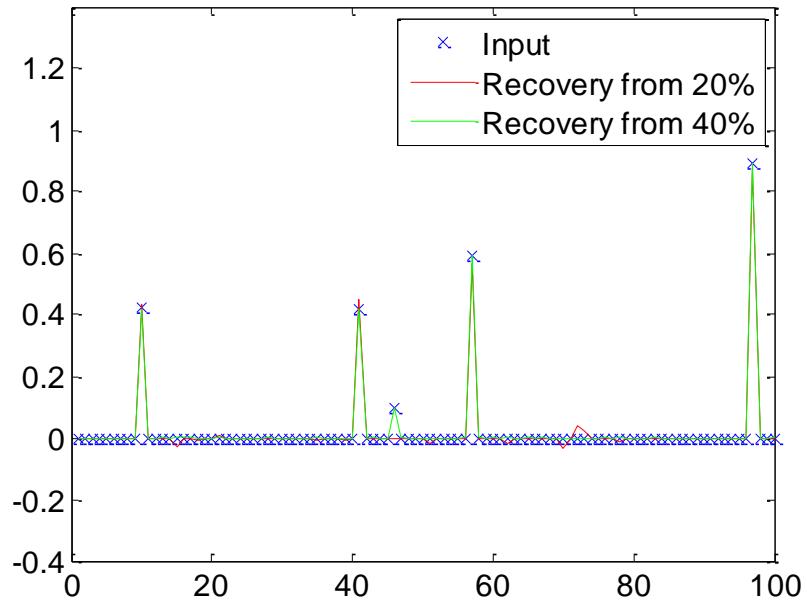
$$t_j = \frac{1}{2}(\mathbb{E}\{X | X \in C_j\} + \mathbb{E}\{X | X \in C_{j+1}\})$$



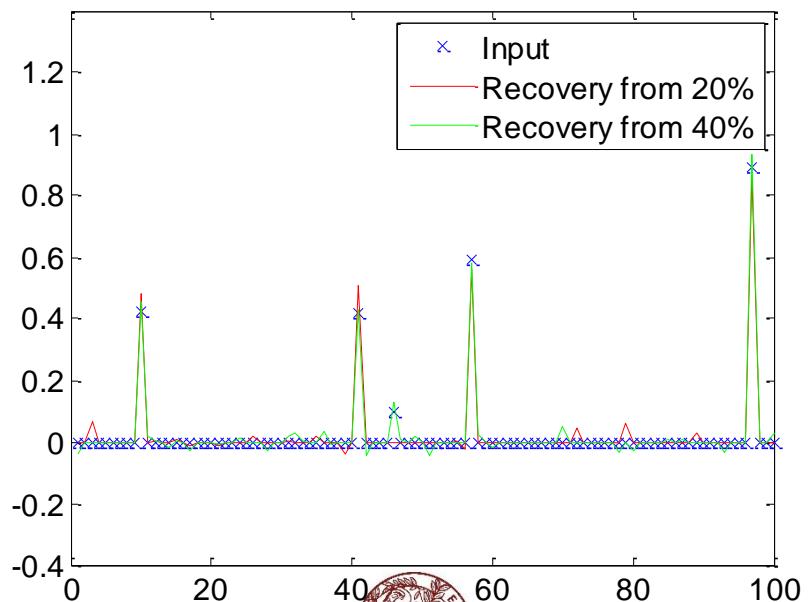
Unquantized Case



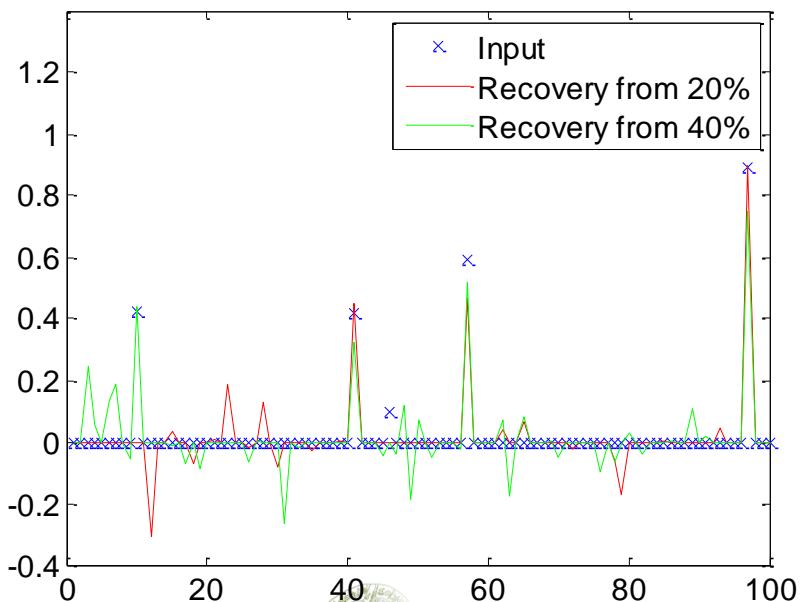
Quantization at 8bpm



Quantization at 4bpm

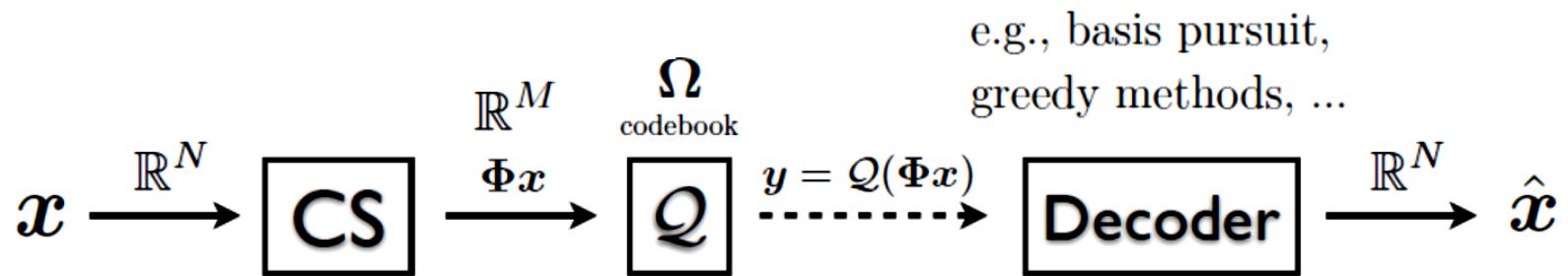


Quantization at 2bpm



Quantizing Compressed Sensing

- In a perfect noiseless system



Finite codebook $\Rightarrow \hat{x} \neq x$

i.e., impossibility to encode continuous domain
in a finite number of elements.

Objective: Minimize $\|\hat{x} - x\|$
given a certain number of:
bits, measurements, or bits/meas.

Addressing Quantization in CS

- Operating regimes

- Measurement Compression
- Quantization Compression

➤ High resolution model (noise source)

$$\Delta \ll \|s\|_2 \longrightarrow \Delta \approx \epsilon_Q : \min \|s\|_1 \text{ s.t. } \|y - \Phi s\|_2 \leq \epsilon_Q$$

➤ 1-bit CS $y = \text{sign}(\Phi s)$

J. Laska, P. Boufounos, M. Davenport, and R. Baraniuk, "Democracy in action: Quantization, saturation, and compressive sensing," *Applied and Computational Harmonic Analysis*, 2011.

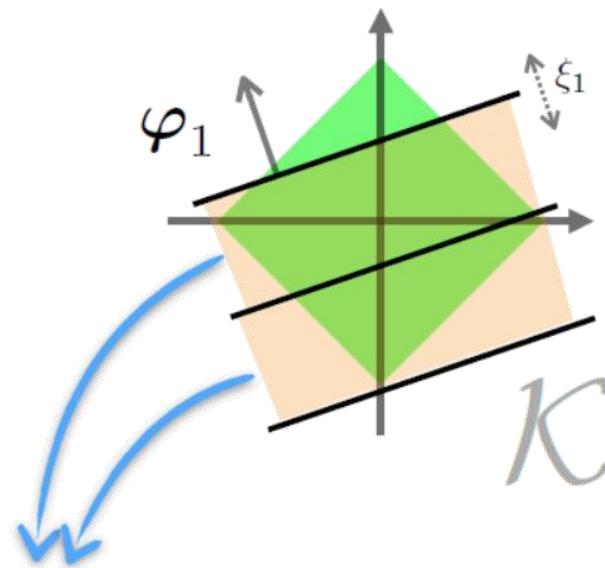
J.N. Laska, and R.G. Baraniuk. "Regime change: Bit-depth versus measurement-rate in compressive sensing." *Signal Processing, IEEE Transactions on Signal Processing*, 2012.

L. Jacques, J.N. Laska, P.T. Boufounos, and R.G. Baraniuk, "Robust 1-bit compressive sensing via binary stable embeddings of sparse vectors," *IEEE Transactions on Information Theory*, 2013.



Properties of $\mathbf{A}(\boldsymbol{x}) := \mathcal{Q}(\Phi\boldsymbol{x} + \boldsymbol{\xi})$

- 1. For consistent reconstruction method



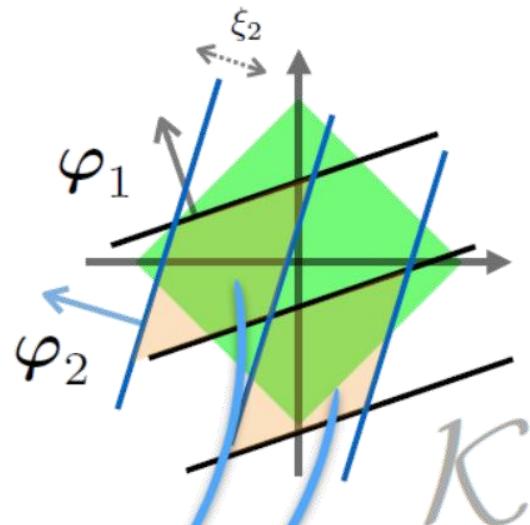
Signals \boldsymbol{u} s.t.

$$\underbrace{\mathcal{Q}(\varphi_1^\top \boldsymbol{u} + \xi_1)}_{\delta[(\varphi_1^\top \boldsymbol{u} + \xi_1)/\delta]} = \text{cst.}$$

$$\Phi = \begin{pmatrix} \varphi_1^T \\ \vdots \\ \varphi_M^T \end{pmatrix}$$

Properties of $\mathbf{A}(\mathbf{x}) := \mathcal{Q}(\Phi\mathbf{x} + \boldsymbol{\xi})$

- 1. For consistent reconstruction method



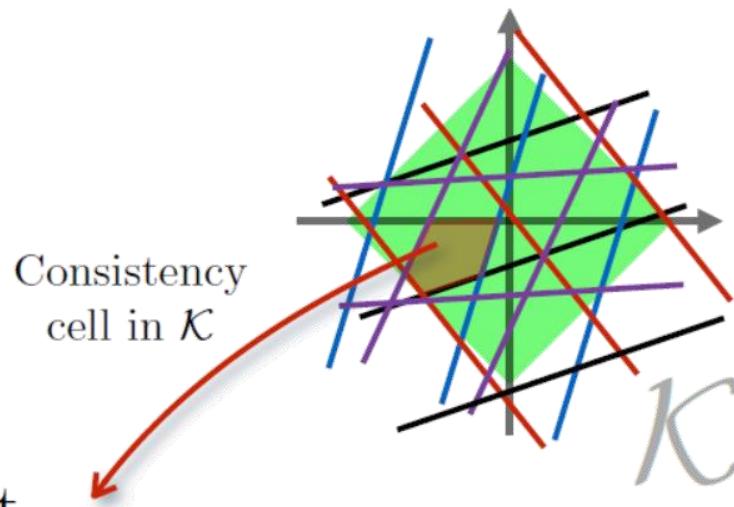
$$\Phi = \begin{pmatrix} \varphi_1^T \\ \vdots \\ \varphi_M^T \end{pmatrix}$$

Signals \mathbf{u} s.t.

$$\left. \begin{array}{l} \mathcal{Q}(\varphi_1^\top \mathbf{u} + \xi_1) = \text{cst.} \\ \mathcal{Q}(\varphi_2^\top \mathbf{u} + \xi_2) = \text{cst.} \end{array} \right\}$$

Properties of $\mathbf{A}(\mathbf{x}) := \mathcal{Q}(\Phi\mathbf{x} + \boldsymbol{\xi})$

- 1. For consistent reconstruction method



Signals \mathbf{u} s.t.

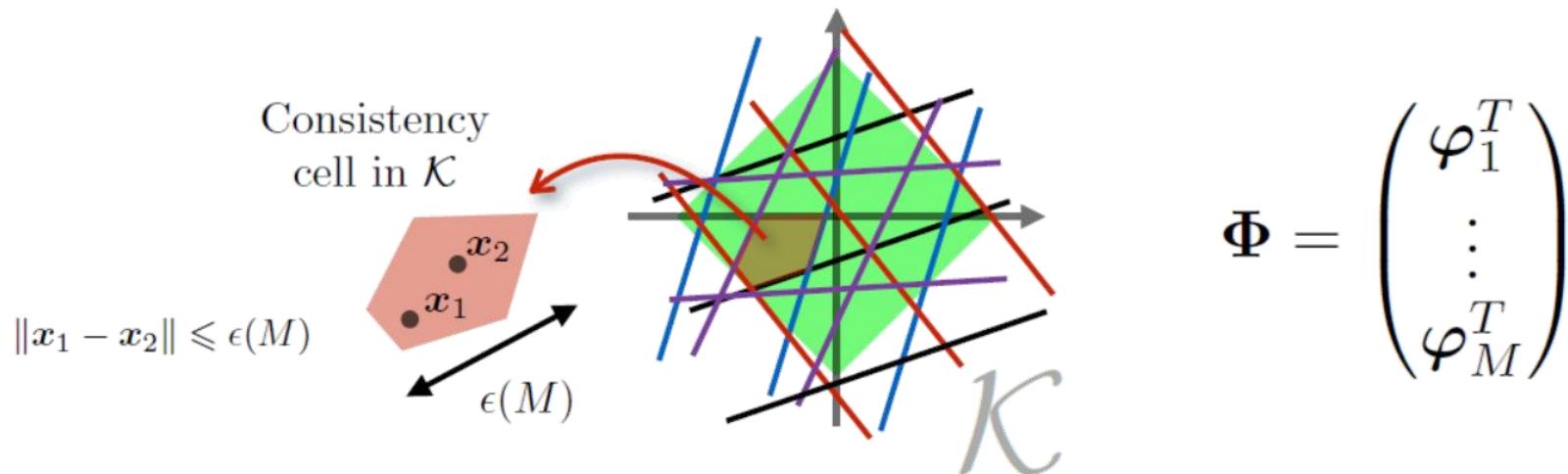
$$\mathbf{A}(\mathbf{u}) := \mathcal{Q}(\Phi\mathbf{u} + \boldsymbol{\xi}) = \mathbf{y}$$

for some \mathbf{y}

$$\Phi = \begin{pmatrix} \varphi_1^T \\ \vdots \\ \varphi_M^T \end{pmatrix}$$

Properties of $\mathbf{A}(\mathbf{x}) := \mathcal{Q}(\Phi\mathbf{x} + \boldsymbol{\xi})$

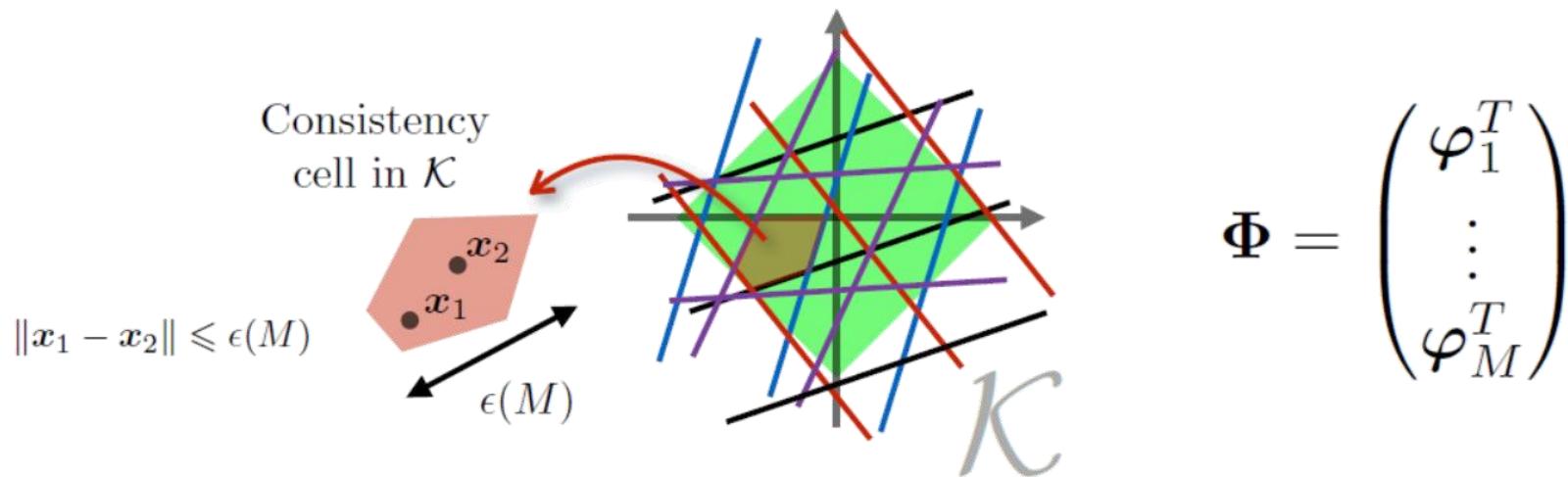
- 1. For consistent reconstruction method



Size $\epsilon(M)$ should decay for large M !

Properties of $\mathbf{A}(\mathbf{x}) := \mathcal{Q}(\Phi\mathbf{x} + \boldsymbol{\xi})$

- 1. For consistent reconstruction method



Size $\epsilon(M)$ should decay for large M !

For Φ a random Gaussian matrix, $\epsilon(M) \leq C_{\mathcal{K}} M^{-q}$.

with $q = 1$ if \mathcal{K} is Σ_K , or low-rank matrices (and others),
and $q = \frac{1}{4}$ otherwise.

[LJ, 17]

Proposed Recovery Mechanism

Introduce *quantization consistency*

Quantized Orthogonal Matching Pursuit (Q-OMP)

- Sparse signals

$$\min \|\mathbf{y} - \mathcal{Q}(\Phi \mathbf{s})\|_2^2 \text{ s.t. } \|\mathbf{s}\|_0 \leq K$$

- Compressible signals

$$\min \|\mathbf{y} - \mathcal{Q}(\Phi \mathbf{Dx})\|_2^2 \text{ s.t. } \|\mathbf{x}\|_0 \leq K$$

Q-OMP

Algorithm 1: Quantized Orthogonal Matching Pursuit (Q-OMP)

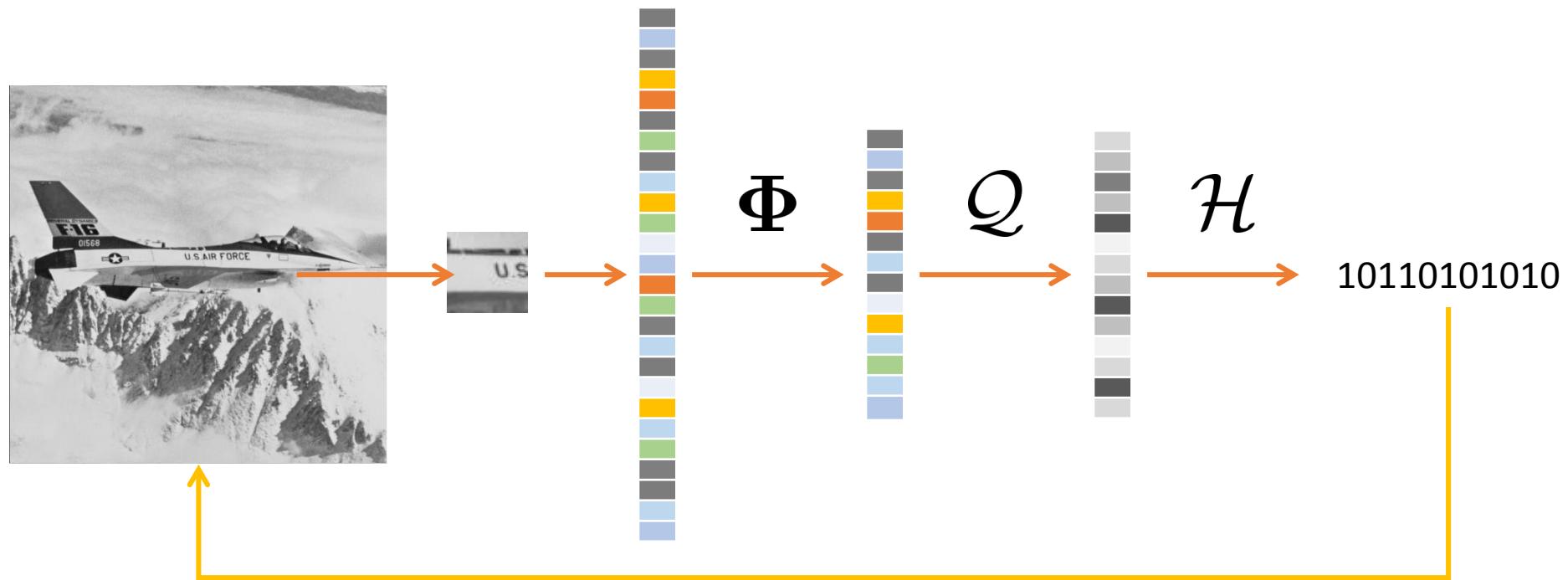
Input: The measurements \mathbf{y} ,
The sensing matrix Φ ,
The dictionary of examples \mathbf{D} ,
The error tolerance *threshold* and/or maximum number of iterations k .
Output: The sparse representation coefficients $\hat{\mathbf{s}}$.

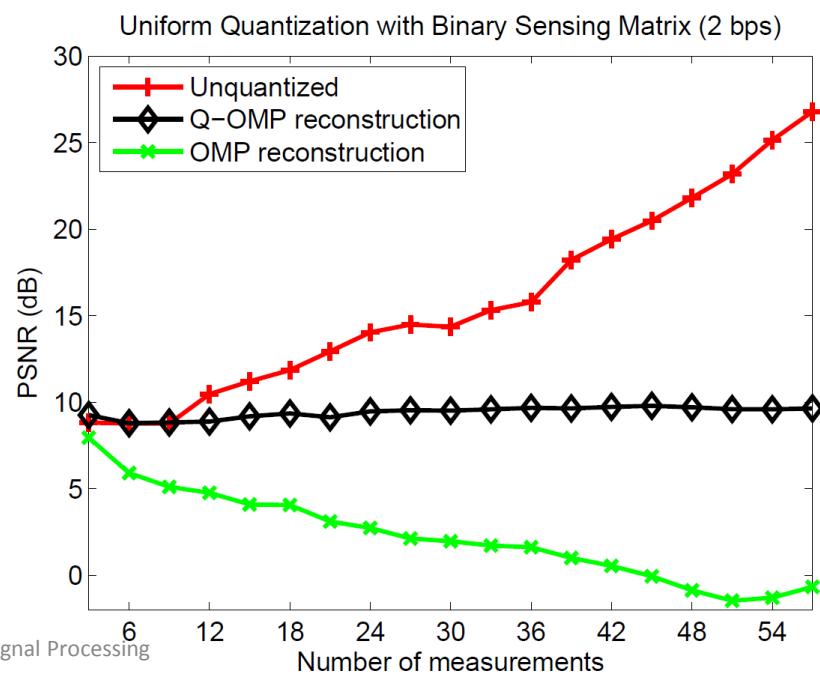
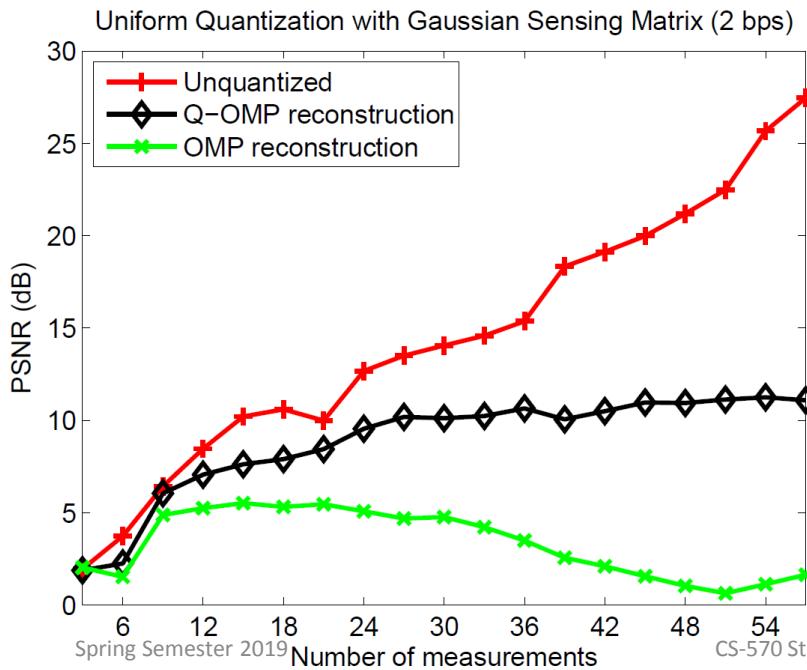
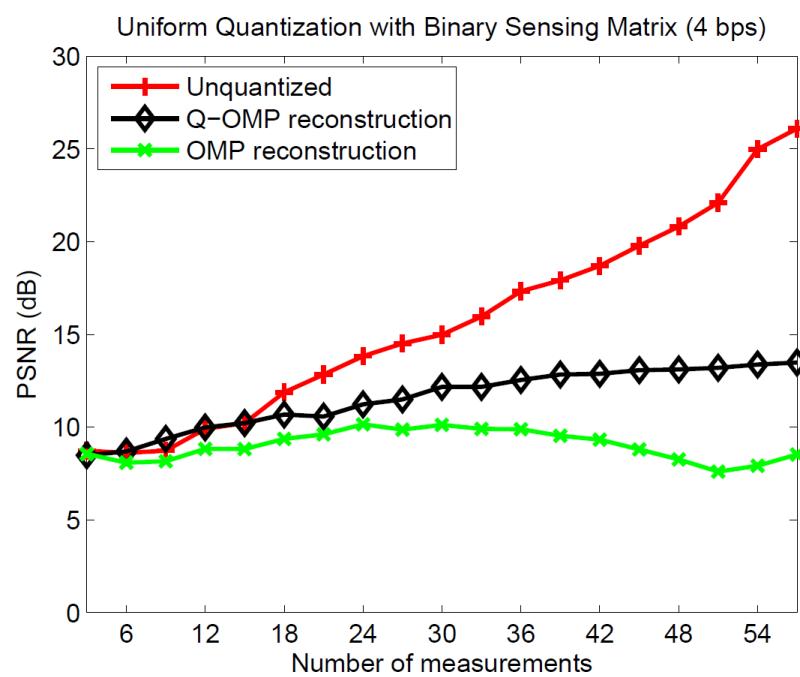
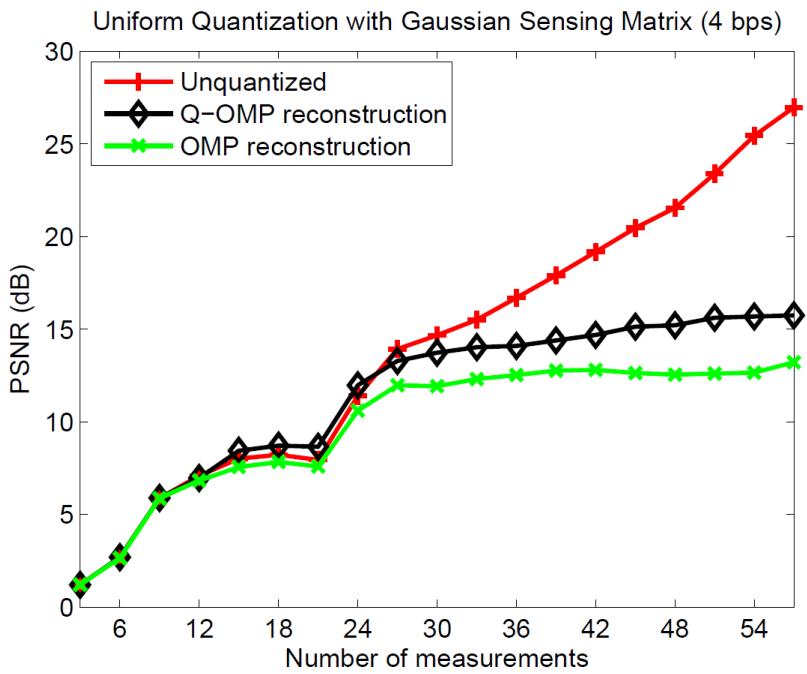
- 1: **initialization** $T^0 = \emptyset, \mathbf{r}^0 = \mathbf{y}$
 - 2: **while** $error \geq threshold$ or $k \leq iteration_lim$ **do**
 - 3: $T^k = T^{k-1} \cup \arg \max_j |\mathcal{Q}(< \mathbf{r}^{k-1}, (\Phi\mathbf{D})_j >) |.$
 - 4: $\hat{\mathbf{s}}_{T^k} = \arg \min_s \|y - \mathcal{Q}((\Phi\mathbf{D})_{T^k}\mathbf{s})\|_2.$
 - 5: $\mathbf{r}^k = \mathbf{y} - (\Phi\mathbf{D})_{T^k}\hat{\mathbf{s}}_{T^k}.$
 - 6: **set** $k \leftarrow k + 1$
 - 7: **end while**
-

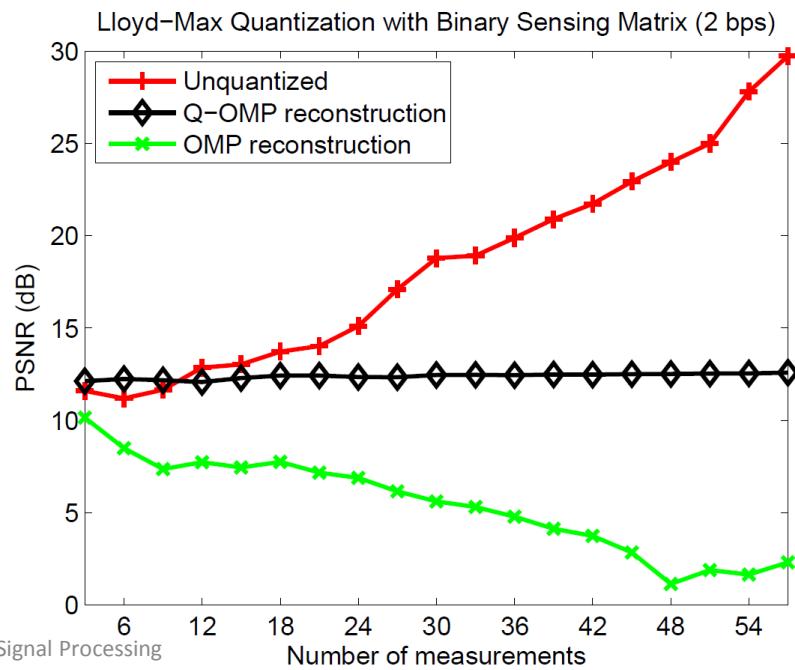
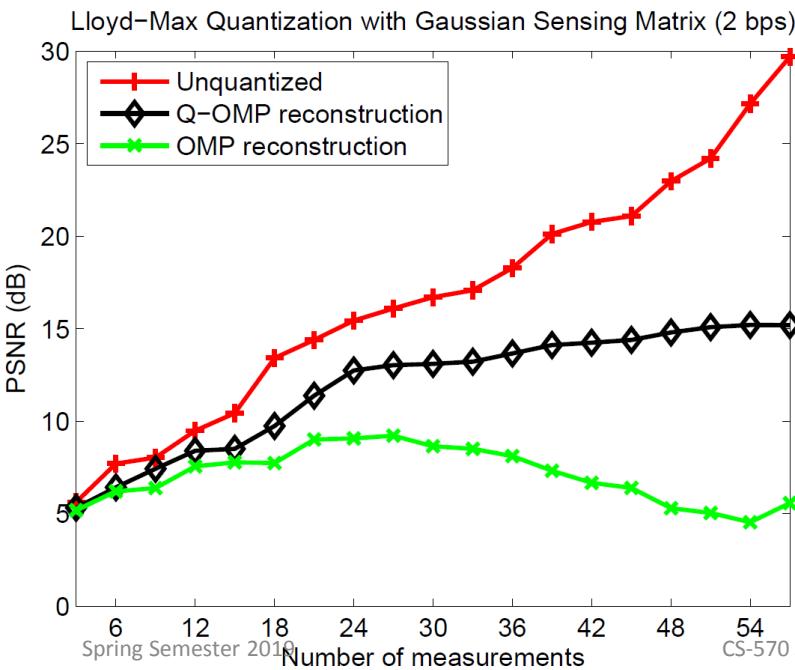
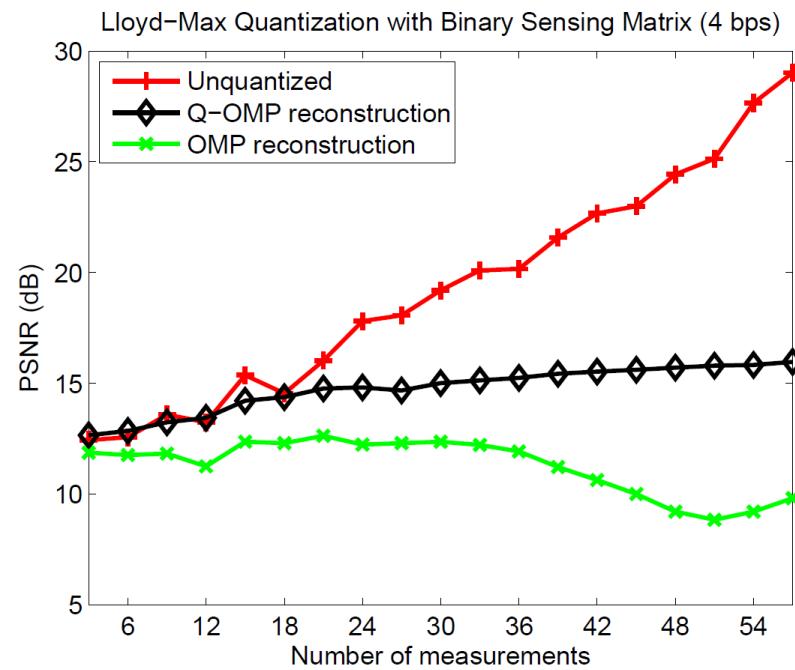
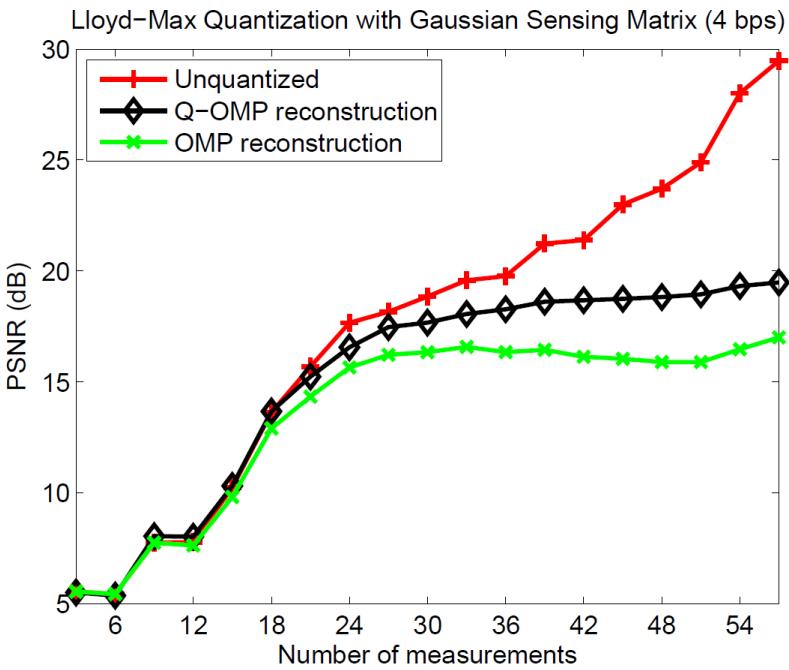


Experimental Setup

Sensing Matrices	Quantization
✓ Gaussian	✓ Uniform
✓ Binary	✓ Optimal







Discussion

		4 BPM	OMP	Q-OMP
Uniform	Gaussian	✓	✓	
	Binary	✗	✓	
Optimal	Gaussian	✓	✓	
	Binary	✗	~	

		2 BPM	OMP	Q-OMP
Uniform	Gaussian	✗	✓	
	Binary	✗	✓	
Optimal	Gaussian	✗	✓	
	Binary	✗		

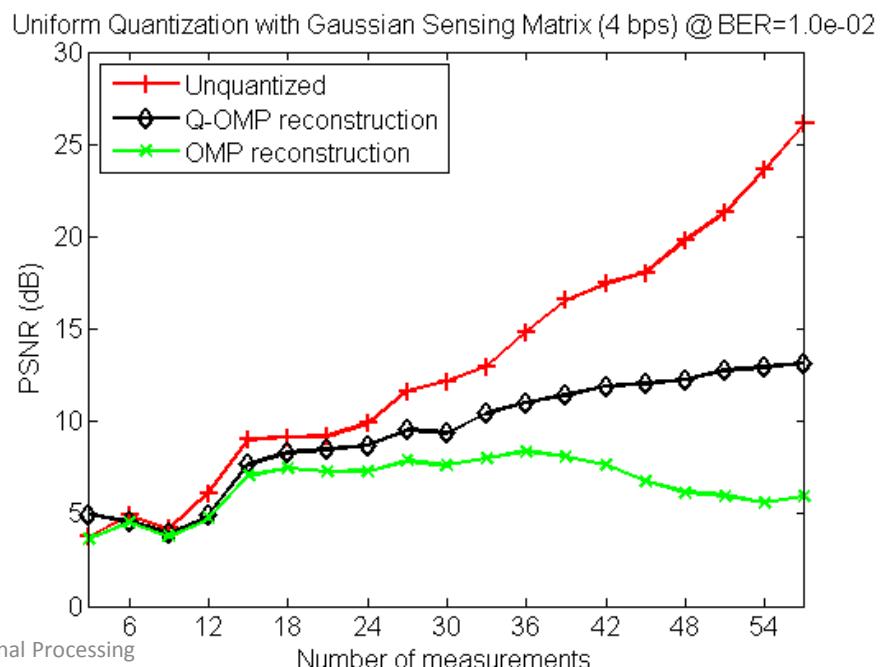
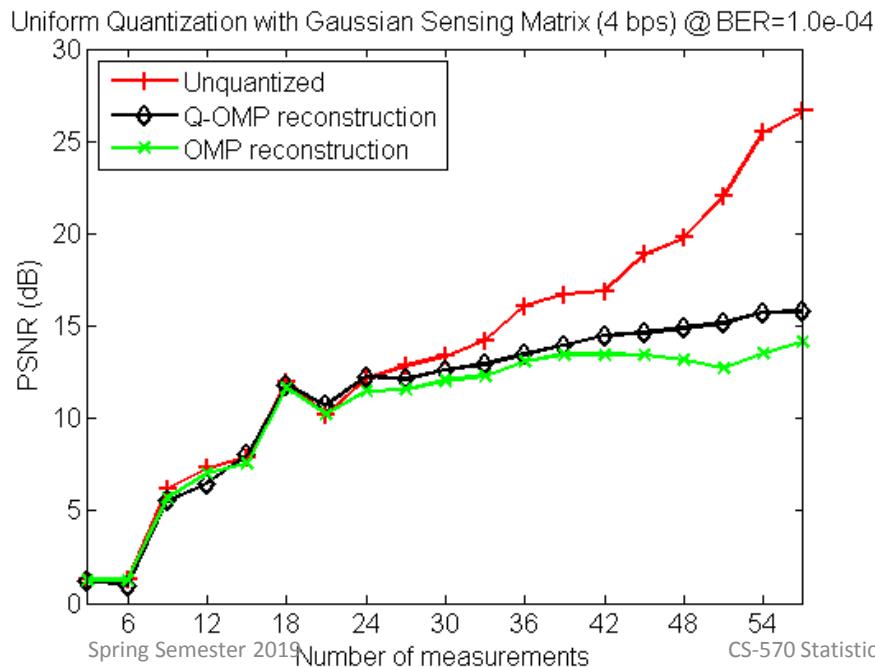
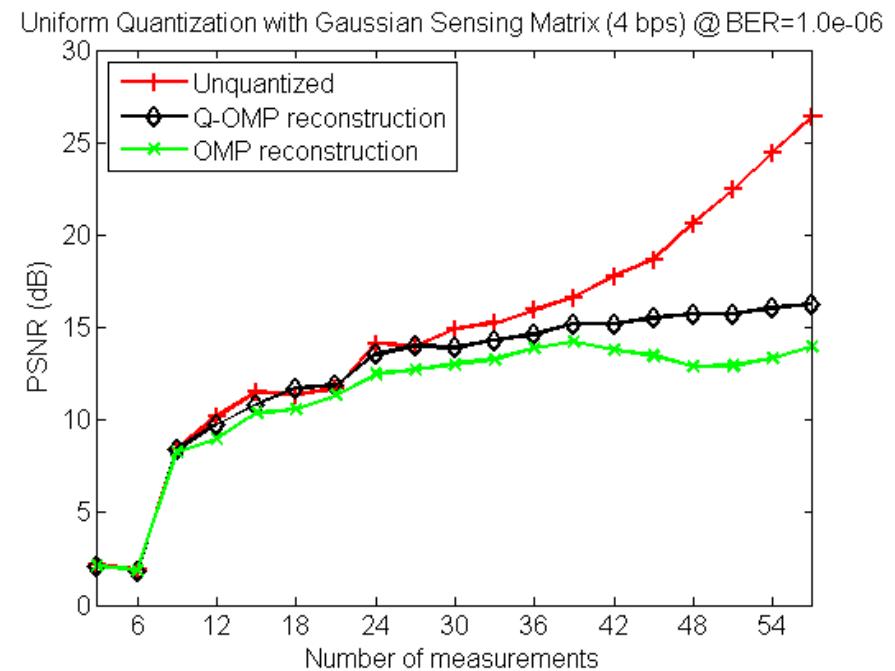
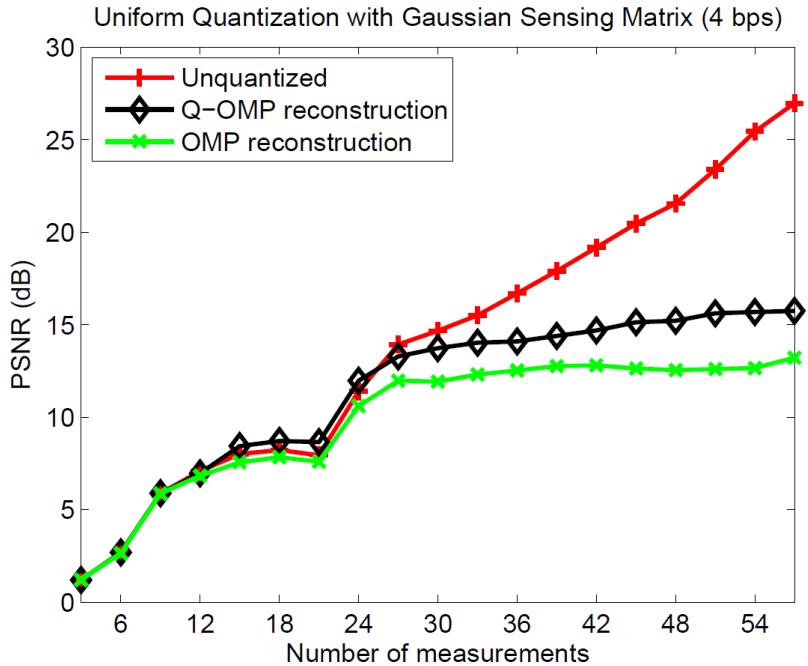
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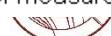
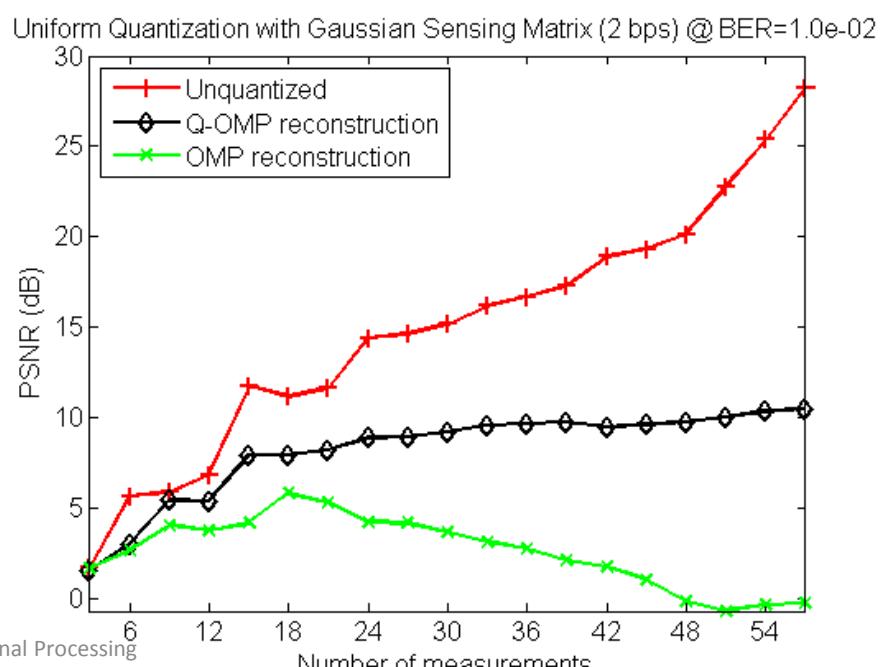
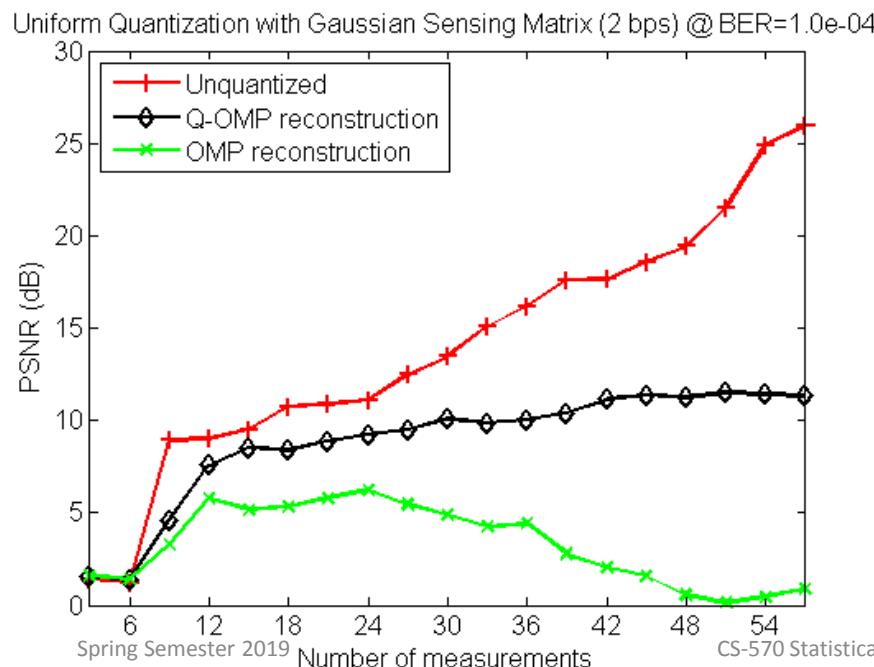
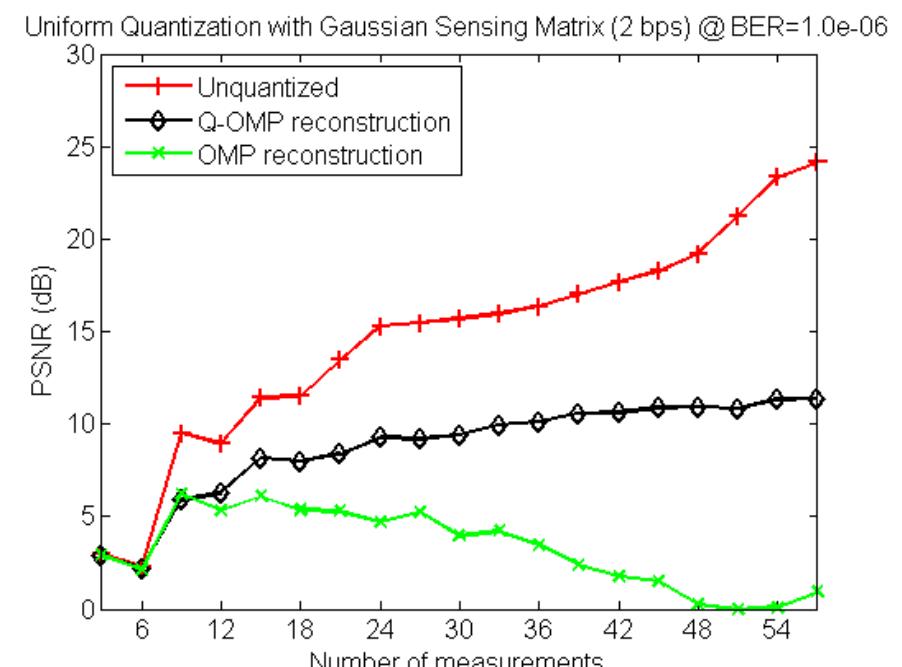
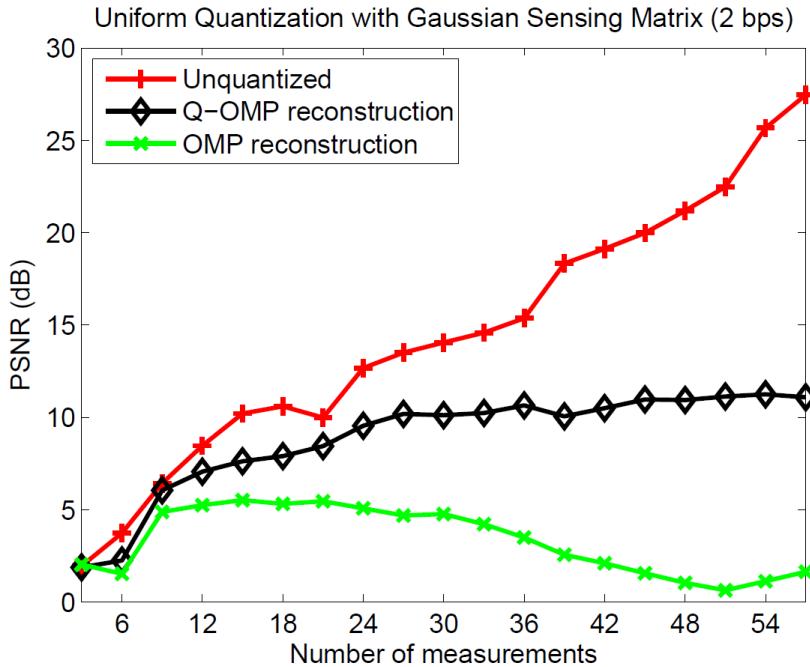


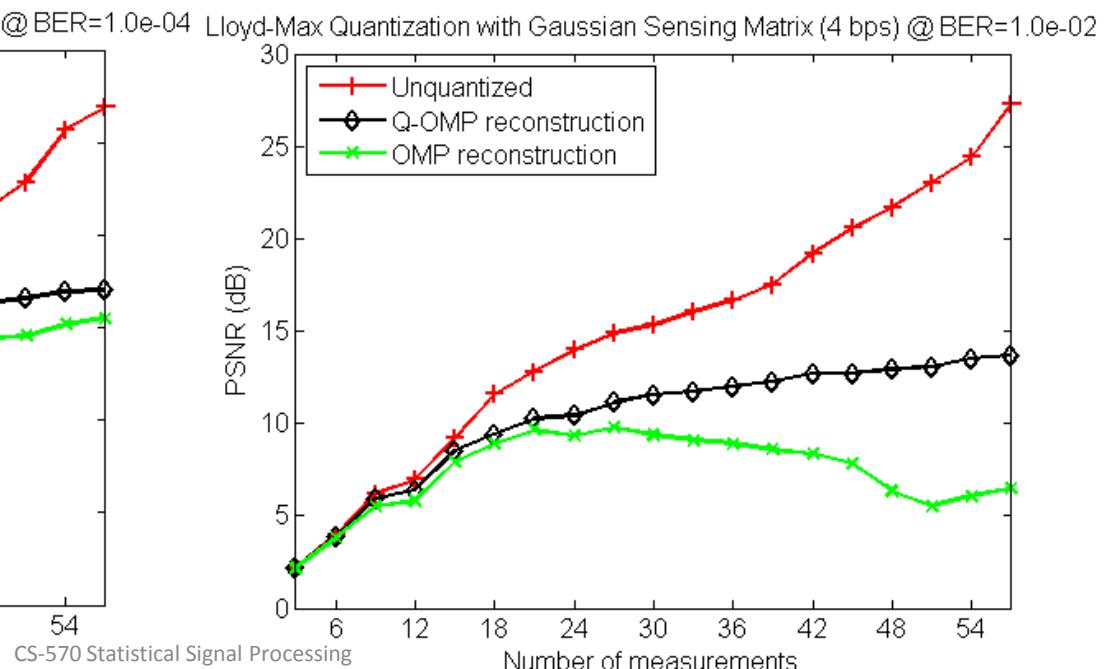
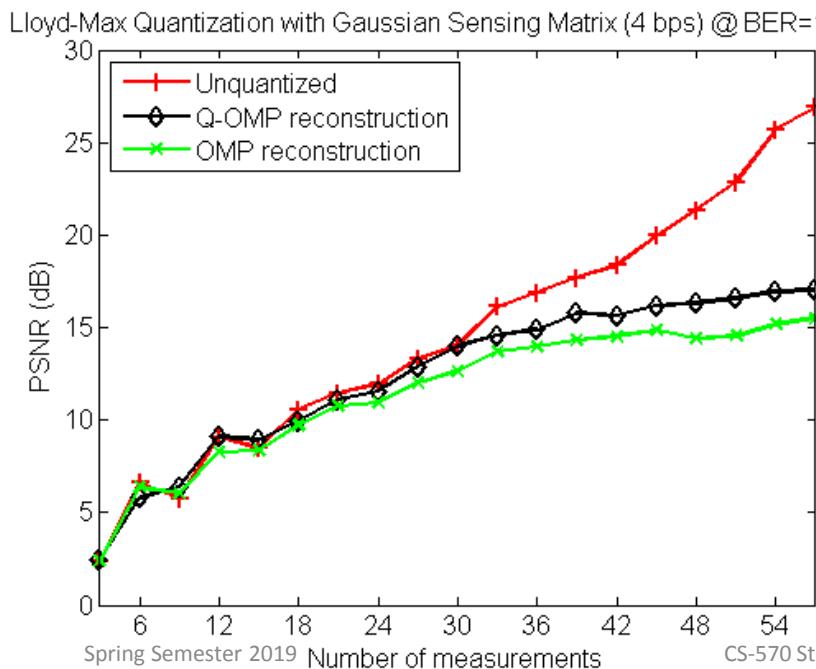
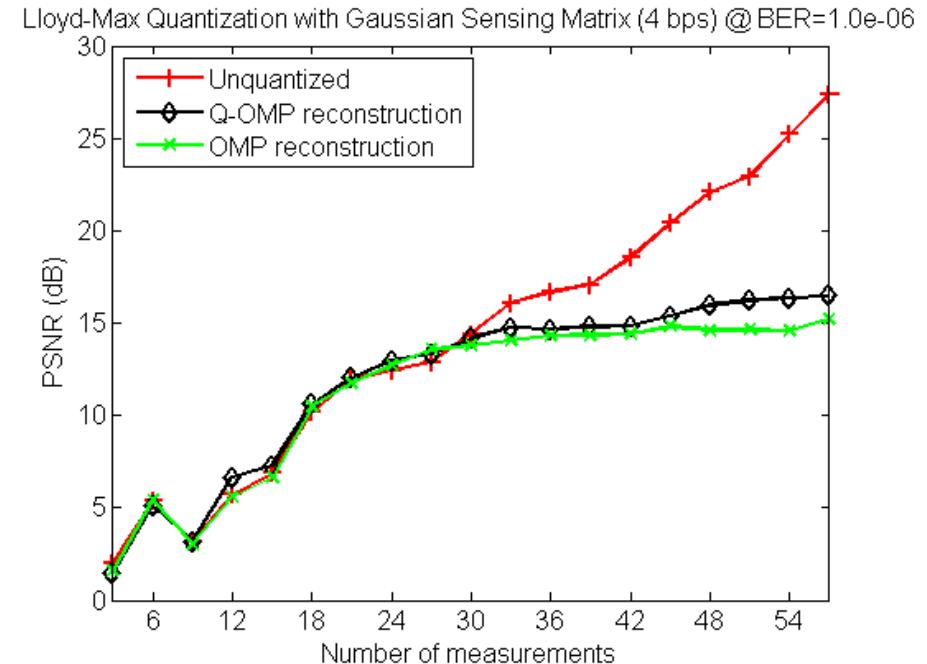
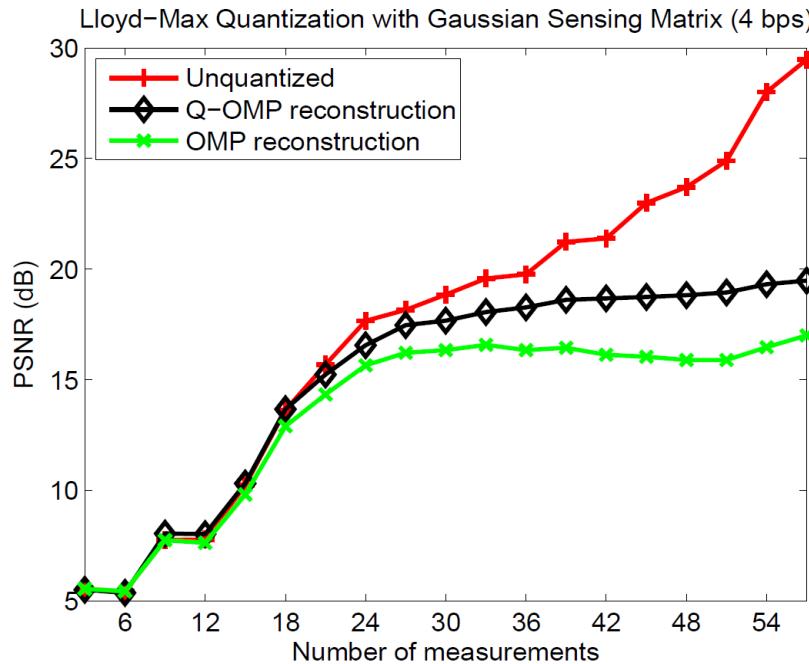
Source coding

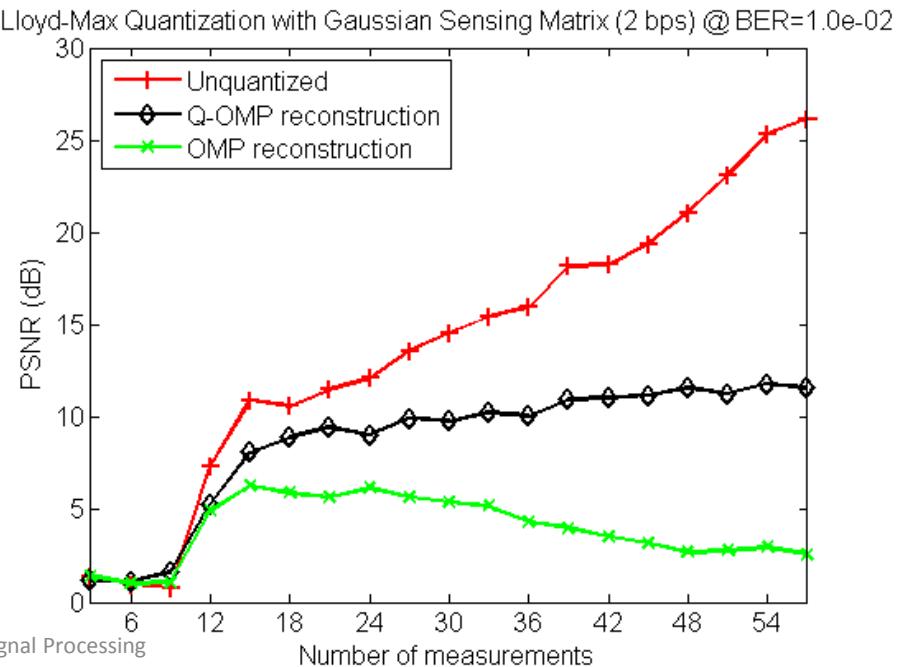
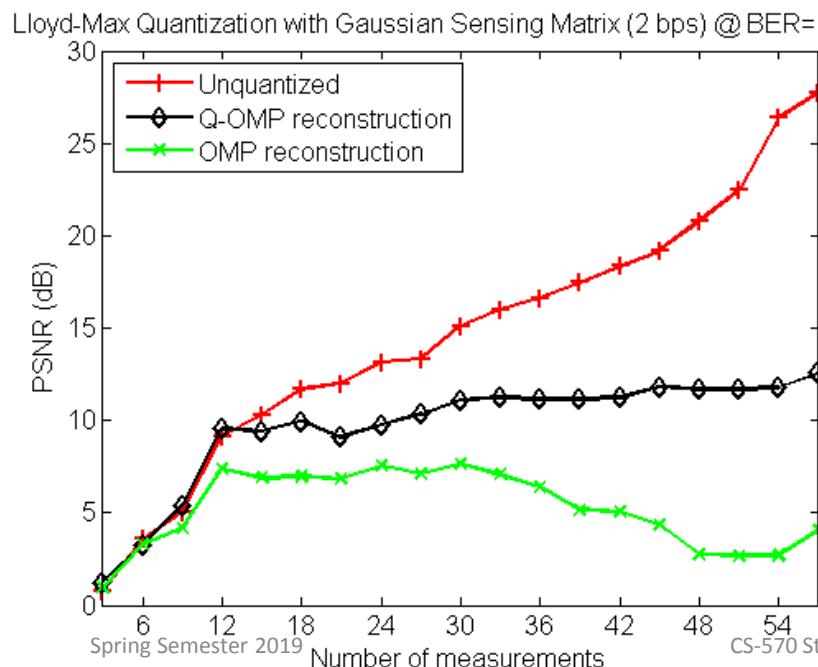
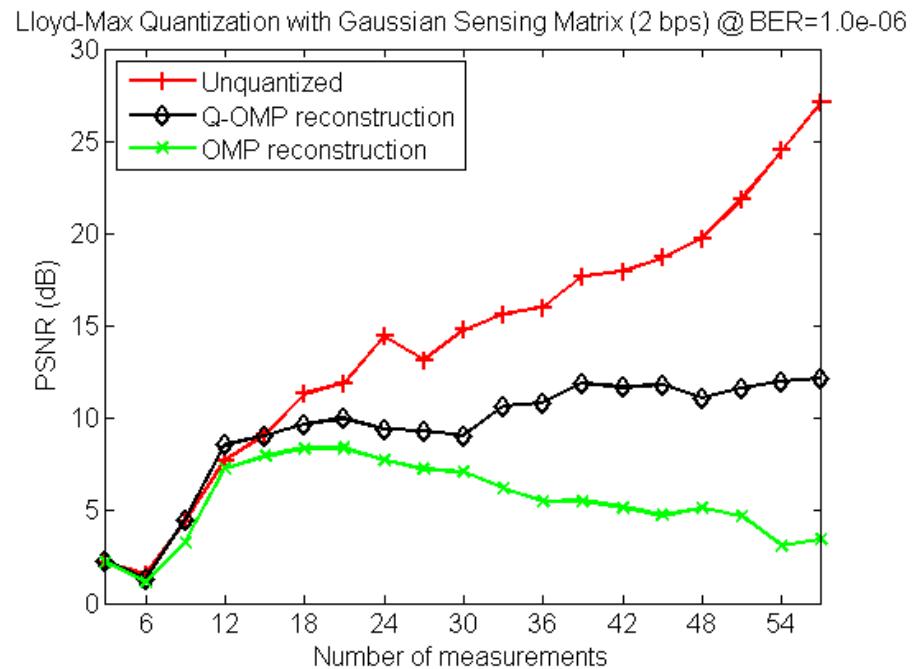
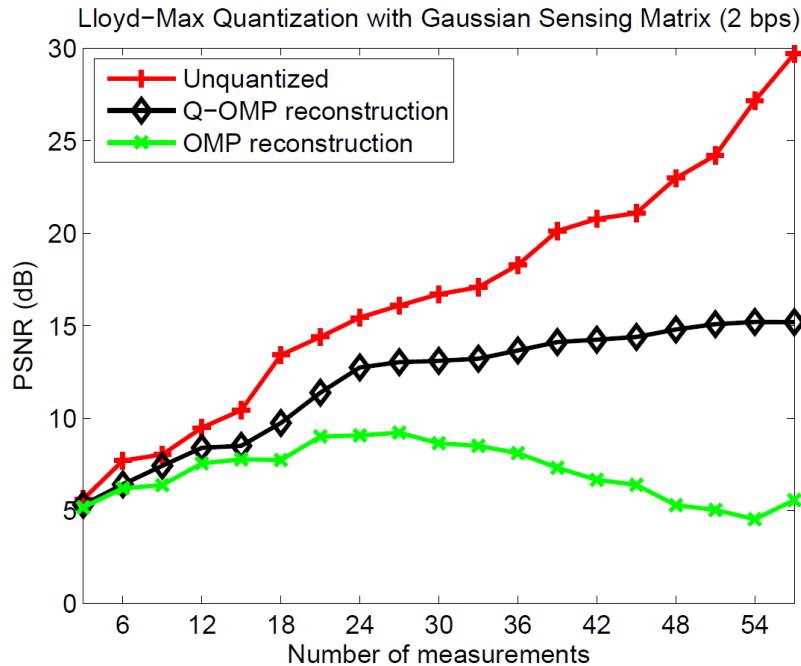
- Examine performance under realistic conditions
- Lossless coding: \mathcal{H}
 - Convert indices to binary vectors
 - Entropy coding
 - Huffman coding(codewords, prob.) -> binary code (JPEG)
 - Arithmetic coding (message) -> binary (JPEG2000)
 - Introduce bit errors
 - Storage
 - Transmission











Discussion

4 b.p.m		OMP	Q-OMP	2 b.p.m		OMP	Q-OMP
Small	Uniform	✓	✓	Small	Uniform	✓	✓
	Optimal	✓	✓		Optimal	✓	✓
Medium	Uniform	✓	✓	Medium	Uniform	✓	✓
	Optimal	✓	✓		Optimal	✓	✓
Large	Uniform	✗	✓	Large	Uniform	✗	✓
	Optimal	✗	✓		Optimal	✓	✓

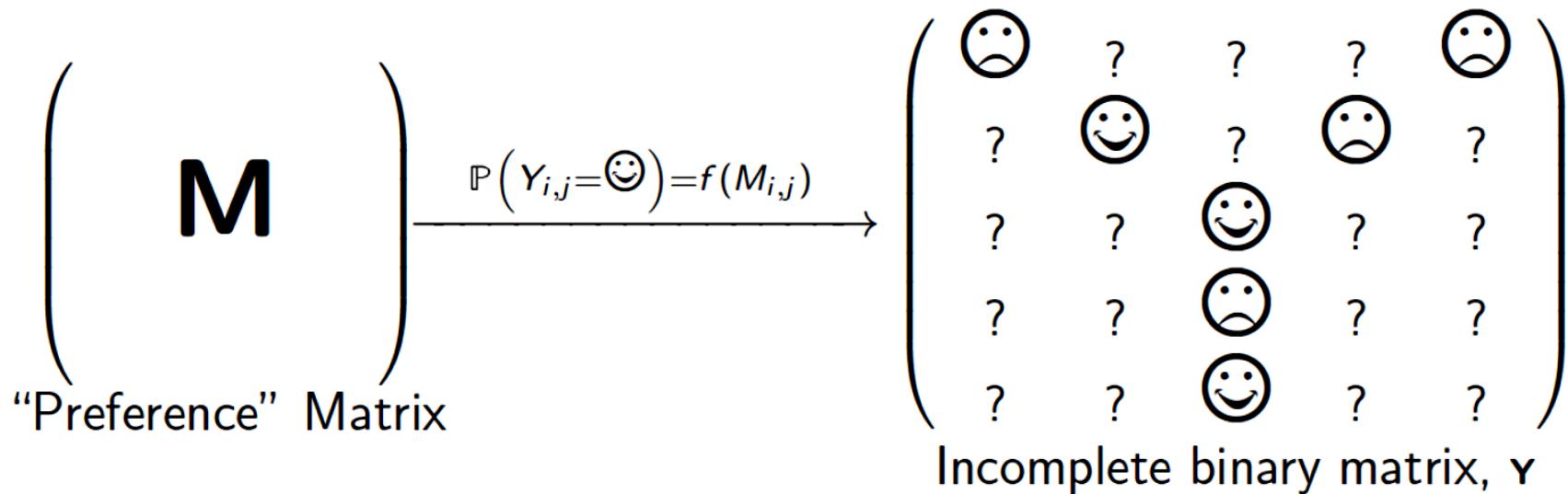
- Quantization is critical in CS
 - Little effect w.r.t uniform vs. scalar
 - Large effects on b.p.m & sampling scheme
 - Sensitive to BER
- Q-OMP
 - Large effect w.r.t optimality
 - Smaller effect on b.p.m & sampling scheme
 - Robust to BER

1-bit Matrix Completion

- Netflix challenge v2

	John	Anne	Scot	Mark	Alice
Chicago	😊	😢	?	?	?
Matrix	😢	?	😊	?	?
Star wars	?	?	😊	?	😢
Inception	?	😊	?	😢	?
Alien	😊	😢	?	?	?
Pulp Fiction	?	?	😢	?	😊

Generalized linear model



- \mathbf{M} is unknown. \mathbf{M} has (approximately) low rank.
- $f : \mathbb{R} \rightarrow [0, 1]$ is a known function (e.g., the logistic curve).
- $\mathbf{M} \in \mathbb{R}^{d \times d}$, $\mathbf{Y} \in \{\odot, \oslash\}^{d \times d}$.
- $\Omega \subset \{1, 2, \dots, d\} \times \{1, 2, \dots, d\}$. You see \mathbf{Y}_Ω .

Formulation

- Consider a $d \times d$ matrix \mathbf{M} with rank r . Suppose we observe a subset Ω of entries of a matrix \mathbf{Y} in the following way:

$$Y_{i,j} = \begin{cases} +1 & \text{if } M_{i,j} + Z_{i,j} \geq 0 \\ -1 & \text{if } M_{i,j} + Z_{i,j} < 0 \end{cases}$$

Observations Matrix Noise

- Generalized model: \mathbf{M} is $d_1 \times d_2$, and $f : \mathbf{R} \rightarrow [0, 1]$

$$Y_{i,j} = \begin{cases} +1 & \text{with probability } f(M_{i,j}), \\ -1 & \text{with probability } 1 - f(M_{i,j}) \end{cases} \quad \text{for } (i, j) \in \Omega.$$

Equivalence when $f(x) = \frac{e^x}{1 + e^x}$ and Z i.i.d logistic distribution



Estimation

- Log-likelihood function

$$\mathcal{L}_{\Omega, Y}(\mathbf{X}) := \sum_{(i,j) \in \Omega} \left(\mathbb{1}_{[Y_{i,j}=1]} \log(f(X_{i,j})) + \mathbb{1}_{[Y_{i,j}=-1]} \log(1 - f(X_{i,j})) \right)$$

- Recovery of \mathbf{M}

$$\widehat{\mathbf{M}} = \arg \max_{\mathbf{X}} \mathcal{L}_{\Omega, Y}(\mathbf{X}) \quad \text{subject to} \quad \|\mathbf{X}\|_* \leq \alpha \sqrt{rd_1d_2} \quad \text{and} \quad \|\mathbf{X}\|_\infty \leq \alpha.$$

- Error

Let f be the logistic function. Assume that $\frac{1}{d} \|\mathbf{M}\|_* \leq \sqrt{r}$. Suppose the sampling set is chosen at random with $\mathbb{E} |\Omega| = m \geq d \log(d)$. Then with high probability,

$$\frac{1}{d^2} \sum_{i,j} d_H^2(f(\hat{M}_{i,j}), f(M_{i,j}))^2 \leq C \min \left(\sqrt{\frac{rd}{m}}, 1 \right).$$

